

4/6 (*) 3.3 Regular Grammar.

$G = (V, P)$ rewriting system.

$V = N \cup \Sigma, N \cap \Sigma = \emptyset, S \in N,$

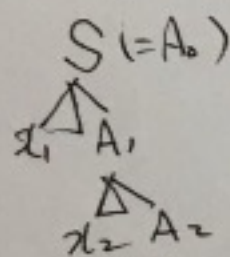
$G = (N, \Sigma, P, S)$ vs $M = (Q, \Sigma, P, q_s, F)$

$P: A \rightarrow x \text{ or } A \rightarrow xB$
 $A, B \in N, x \in \Sigma^*$

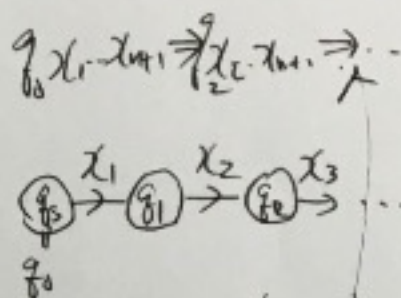
$P: q \cdot x \rightarrow p \in P$ vs $p \in \delta(q, x)$
 $p, q \in Q, x \in \Sigma^*$

right linear grammar (type 3)

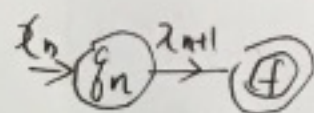
left linear grammar
 $A \rightarrow x \text{ or } A \rightarrow xB$



1:1



$\Rightarrow q_n \cdot x_{n+1} \Rightarrow F$



terminal string consumed

linear structure

vs context free grammar (type 2)
 hierarchical structure

skewed tree \equiv linear list

$S \Rightarrow x_1 A_1 \Rightarrow x_1 x_2 A_2 \Rightarrow \dots \Rightarrow x_1 \dots x_n A_{n+1} \Rightarrow x_1 \dots x_n x_{n+1}$

generating

* program

non-hierarchical structure
 - multiple inheritance

vs single "tree" graph

global variable considered harmful!

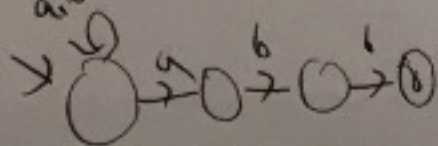
* $\frac{1}{1+0.01}$ - return value

multiple-value returning function!

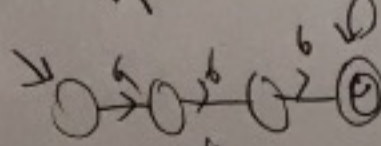
3.4 Deterministic FA

$q \cdot w \Rightarrow q_1 w_1$
 $q \cdot w \Rightarrow q_2 w_2$
 $r_1 \neq r_2$ non-deterministic

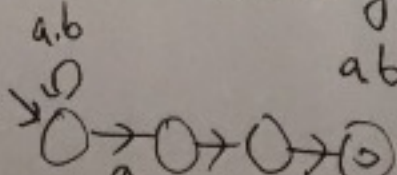
normal form $|x|=1$ $|p|=1$ $x=0$ - unique transition



suffix



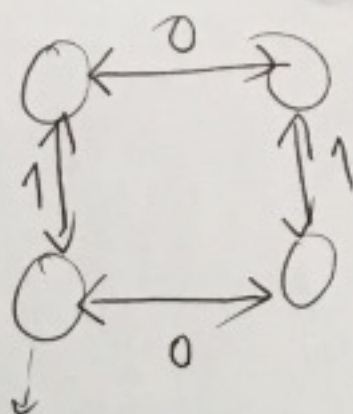
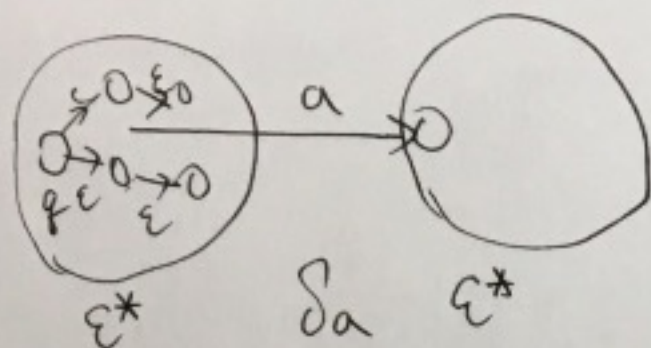
prefix



substring

③ ϵ -NFA \Rightarrow DFA

DFA



ϵ -closure

$$\delta_a^* \equiv \delta_\epsilon^* \cdot \delta_a \cdot \delta_\epsilon^*$$

vs $\hat{\delta}_a \equiv \delta_\epsilon^* \cdot \delta_a$

$$[0110111] = [1] = [0^{\frac{22}{7}} | \frac{1}{2}] =$$

$$[001] = [1] = [100] = [010]$$

① NFA \rightarrow DFA : subset construction