第16节 6.1. Viable Prefix = viable stack string.

Gab:

$S \to aA \mid bB$

$A \to c \mid dAd$

$B \to c \mid dBd$

$L(G) = \{ t^n c d^n \mid n \geq 0 \}$ $\overset{\alpha}{=}$ $|a \to a|$ $|\alpha| \to A|$

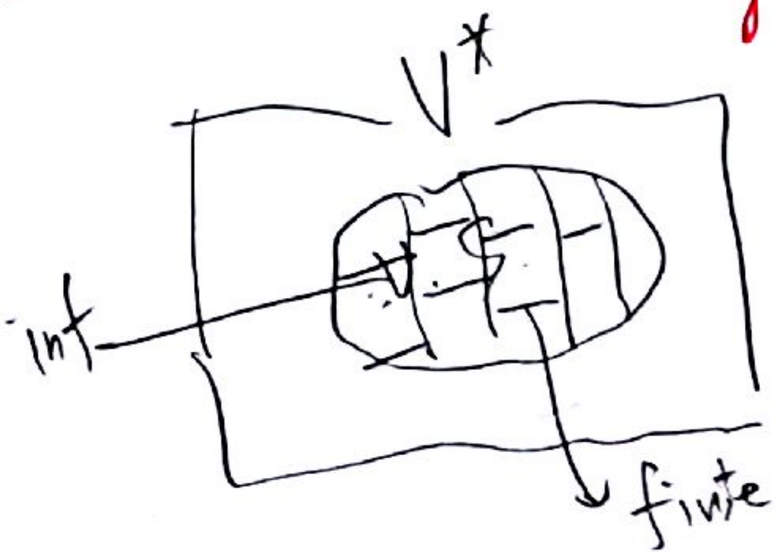
$\{a,b\}^*$

$\alpha c \mid x \to \alpha A \mid x$

$\beta c \mid y \to \beta B \mid y$

lookback lookahead

pef. viable stack strig



V.S. of G_{ab} ad^* ad^*c reduce c to A

$\varepsilon \{\varepsilon\} \cup \{ad^n \mid n \geq 0\} \cup \{ad^n c \mid n \geq 0\}$

$\cup \{ad^n A \mid n \geq 0\} \cup \{ad^n Ad \mid n \geq 1\}$

$\cup \{b \overset{ad^*A}{\cdots}\}$ $\overset{odd^*Ad}{\underset{d^*}{\longrightarrow}}$ reduce c to B

$\cup \{S\}$

valid actions for viable stack string γ_1, γ_2

same set of two

$|\Sigma| + |P|$ actions $2^{|\Sigma|+|P|}$

shift reduce

~~shift~~ right parser

$a \in \Sigma$ shift a,

$A \to X_1 X_2 \cdots X_n \in P$ reduce $X_1 \cdots X_n$ to A,

$|a \to a|$

$|X_1 X_2 \cdots X_n| \to A|$

shift-reduce parser

$[\delta] \mid a \to [\delta][\delta a] \mid$

$[\delta][\delta X_1][\delta X_1 X_2] \cdots [\delta X_1 X_2 \cdots X_n] \mid \to [\delta][\delta A]$