제 21강 (5/27 화) LR parsing.

왼 파서 Top-down

Left parser $L = (\quad ; \quad)$ $\underline{L(L) = L(G)}$

(produce-shift parser)

$A| \rightarrow \alpha^R|$ ... produce, guess A as α $\quad A \rightarrow \alpha \in P$

Top

$a|a \rightarrow |$ shift, verify a $\quad a \in \Sigma$

$A \to \alpha | \beta$ guess-guess conflic

Non det

$\alpha = X_1 X_2 \dots X_n$

$A|X_n \cdots X_2 X_1$

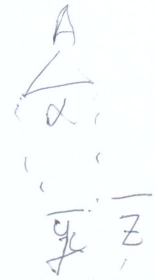
α^R

SLL(k) parser $SLL_k = (\quad)$

$A|x \rightarrow \alpha^R|x$ $\quad A \to \alpha \in P$ $\quad x \in \underbrace{First_k(\alpha)}_{y} \oplus \underbrace{Follow_k(A)}_{z}$

adding lookahead

$a|a \rightarrow |$ $\quad a \in \Sigma$



$k : yz$

$\underline{L(A) \supseteq L(SLL_k)}$
$\subseteq$

오른 파서

Right Parser Bottom-Up $\quad R \quad L(R) = L(G)$

(shift-reduce parser)

$|a \rightarrow a|$ shift $- a \in \Sigma$

$\alpha| \rightarrow A|$ reduce α to A, $A \to \alpha \in P$

stack symbol $X \in N \cup \Sigma$

" string $\gamma \in (N \cup \Sigma)^*$

" $\gamma X \rightarrow [\gamma X]_R$

연역 理 · 學而之 銘名 – 논리 공자의 패러다

이 귀납 氣 · 名可名, 非常名 – 노자 도덕경

기 명가명, 비상명

$R \subseteq (N \cup \Sigma)^* \times (N \cup \Sigma)^*$

equivalence class

equiv relation

$\gamma_1 R \gamma_2$

if ~~~

stack strg $(N \cup \Sigma)^*$ – infinite

LR(k) state partition

partition – finite