

11/22(화) 제21강 Deterministic Right Parsers

Review Deterministic Left Parsers.

$$\begin{aligned}
 E &\rightarrow E+T \mid T \\
 T &\rightarrow T * F \mid F \\
 F &\rightarrow a \mid (\epsilon)
 \end{aligned}$$

$+$: left asso.
 $*$: " "
 $\text{prec}(* > \text{prec}(+)$

$$\begin{aligned}
 E &\rightarrow T+E \mid F+T \mid a \mid (\epsilon) \\
 T &\rightarrow F * T \mid a \mid (\epsilon) \\
 F &\rightarrow a \mid (\epsilon)
 \end{aligned}$$

$+$: right asso.
 $*$: " "
 $\text{prec}(* > \text{prec}(+)$

TP p6.

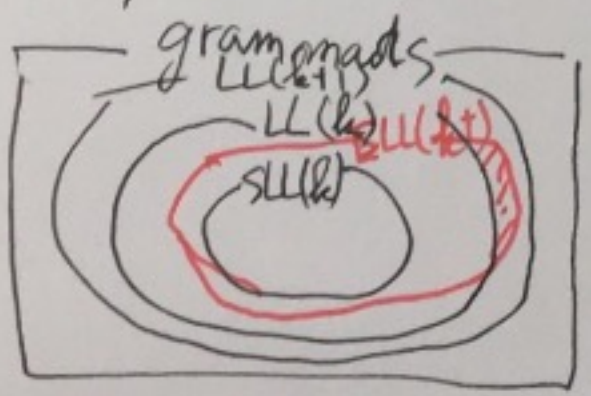
$$\begin{aligned}
 T: \text{First}_1(F T') &= \{a, (\epsilon)\} \\
 E: \text{First}_1(F T E') &= \{a, (\epsilon)\} \\
 E': \text{First}_1(T E') &= \{a, (\epsilon)\}
 \end{aligned}$$

SLL(k) parser

$$A \rightarrow \alpha \in P \quad x \in \text{First}_k(\alpha) \cap \text{Follow}(A)$$

$$(x, \cancel{A}) \xrightarrow{A \rightarrow \alpha} (x, \alpha) \quad \dots \text{SLL(k) parser} \dots$$

$$(\epsilon, A) \xrightarrow{A \rightarrow \alpha} (\epsilon, \alpha) \quad \dots \text{Left parser} \dots \text{non det.}$$



→ Deterministic Right Parsers.

LR(0) item $[A \rightarrow \alpha \cdot \beta]$ if $A \rightarrow \alpha \beta \in P$.

* Size of grammar G

$$|G| = \sum_{A \rightarrow \alpha \in P} |A| + |\alpha|$$

~~the~~ D. Knuth (1965)

On the Translation of Languages
from Left to Right.

$$LR(k) = LR(1)$$

$$\supseteq LALR(1)$$

$$\supseteq SLR(1)$$

$$\supseteq LR(0)$$