

총 28강 중 27강 남음

Review: Chap. 7. Useful grammar, $\xrightarrow{\text{X}\epsilon\text{-rule, X unit}} \text{Chomsky's Normal Form}$
 $A \rightarrow BC \vee a$ (CNF for cfg)
 except $S \rightarrow \epsilon$.

CFG $A \rightarrow \alpha \in P, A \in N, \alpha \in (N \cup T)^*$

\hookrightarrow normal form $A \rightarrow BC \vee A \rightarrow a, A, B, C \in N, a \in T$
 except $S \rightarrow \epsilon$.

RG. $A \rightarrow xB \vee A \rightarrow y \in P, A, B \in N, x, y \in T^*$

\hookrightarrow normal form $A \rightarrow aB \vee b \in P, A, B \in N, a, b \in T$. \uparrow Deterministic Reg. grammar

DFA $\delta: Q \times \Sigma \rightarrow 2^Q$ (or D) $\xrightarrow{\text{If}} A \rightarrow xB, A \rightarrow aC \in P$, then $B=C$.

XFA $\delta: Q \times \Sigma^* \rightarrow 2^Q$

$\text{If } B \neq C, \text{ then } A \rightarrow aB, aC \notin P$.

Pumping Lemma for CFL

$w = uvwx^i y$

- i) First pump $|vwx| \leq n$
- ii) non-empty pumping $vx \neq \epsilon$
- iii) pumping $uv^iwx^i y \in L, \forall i \geq 0$.

Pumping Lemma for RL

$w = xyz$

- i) $|xy| \leq n$
- ii) $y \neq \epsilon$
- iii) $x(y^i)z \in L, \forall i \geq 0$.

~~CFL~~ Closure properties of CFL's
 union, Concatenation, closure, Substitution, Reversal

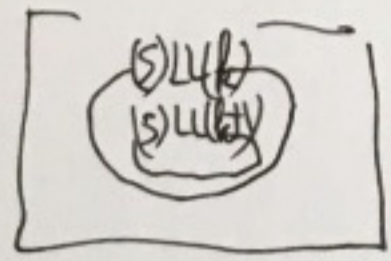
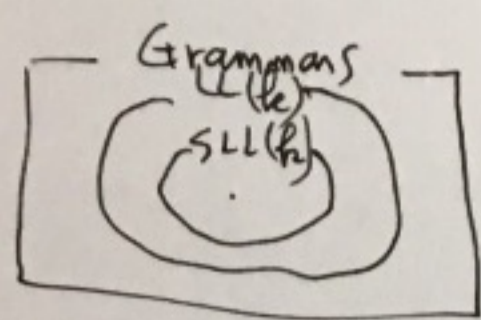
\times intersection But $CFL \cap RL \Rightarrow CFL$

Membership Problem

CYK $O(n^3)$
 left/right NP $(O(2^n)!)!$

3) chap 6-B Deterministic Left Parser.

SLL(k) ... Strong Left-to-right scan in Leftmost derivations with k-lookahead symbols.
LL(k)



But $SLL(1) = LL(1)$.

Left Parser

$(\epsilon, A) \rightarrow (\epsilon, \alpha) \quad A \rightarrow \alpha \in P \quad (a, a) \rightarrow (\epsilon, \epsilon) \quad a \in T$

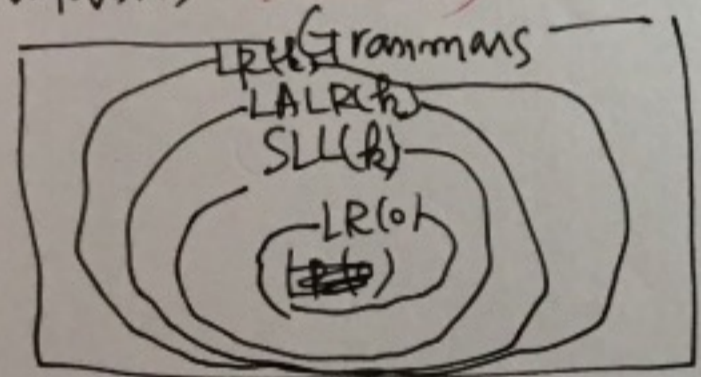
SLL(k) Parser

$(x, A) \rightarrow (A, \alpha) \quad A \rightarrow \alpha \in P \quad (a, a) \rightarrow (\epsilon, \epsilon) \quad a \in T$

- i) $x \in First_k(\alpha)$
- ii) Left recursion \rightarrow right recursion

ii) $x \in First_k(\alpha) \oplus_k Follow(A) \quad A \rightarrow \alpha \in P$

Right Parsers SLL(k)



$LR(0) \subsetneq SLL(k) \subsetneq LALR(k) \subsetneq LR(1)$

$SLL(k) \subset SLL(k+1)$

$LALR(k) \subset LALR(k+1)$

But $LR(1) = LR(k) \dots$ Knuth, 1965

"One the translation of the languages from Left to Right"