

2/21/20 11/29 (EU) Turing-Church's Thesis.

Review

TM의 계산/결정

GTT (Gödel's incompleteness Theorem) - Denial of Self Recursion.

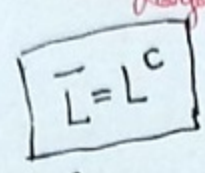
Partial recursive function
primitive recursive function
 μ -recursive function

Complement of Recursive & recursively enumerable languages

Def. R.E. lang $L \leftrightarrow \exists a \text{ TM } M \forall x \in \Sigma^* \langle x \in L \rightarrow M \text{ accept} \rangle$
 $\langle x \notin L \rightarrow M \text{ reject or does not terminate} \rangle$

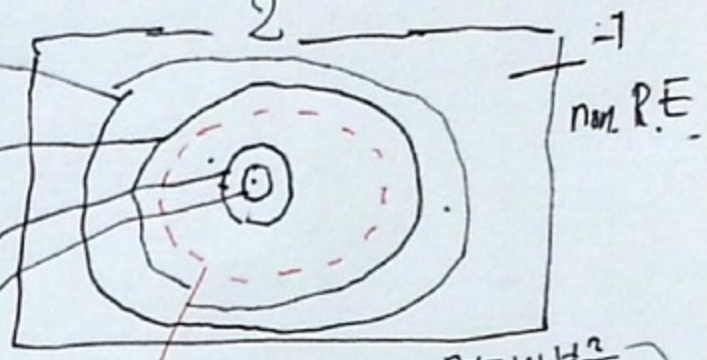
recursive lang. $L \leftrightarrow \exists a \text{ TM } M \forall x \in \Sigma^* \langle M \text{ terminate} \rangle$

Fact $\exists a \text{ TM } M \exists L = L(M)$
 $\Rightarrow \exists a \text{ TM } M' \exists \bar{L} = L(M')$



$L(M) = \{x \in \Sigma^* \mid (\exists \alpha, \beta, \gamma) \vdash^* \langle \alpha, f, \beta \rangle \Rightarrow x, \beta \in \Gamma^*, f \in \Gamma\}$

$\bar{L} = U - L$
 $L = \{x \in U \mid p(x)\}$
 type 3. 정지 문제
 fa



NP-complete

Thm-증명	Thesis
Lemma	아직 증명 X
Coll. 사례	
Fact (명제)	

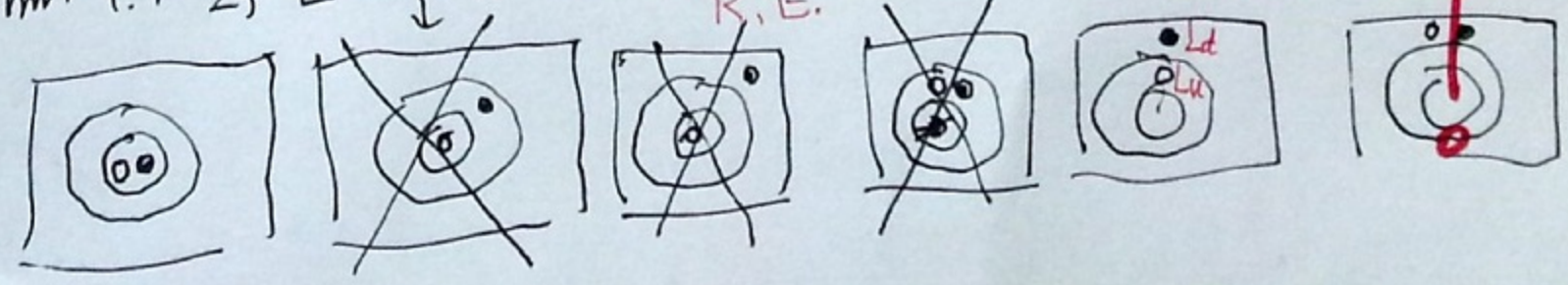
configuration of TM.
 (I.D.) $P^* \times Q \times P^*$ $P: \text{일} \neq \text{일}, \text{memory}$

Configuration of PDA
 $Q \times P^*$ $P: \text{일} \neq \text{일}, \text{stack}$

conf. of FA
 $Q \times \Sigma^*$ $\Sigma: \text{일} \neq \text{일}$

Thm 9.3 If L is recursive. Then \bar{L} is also
 accept \rightarrow reject
 reject \rightarrow accept.

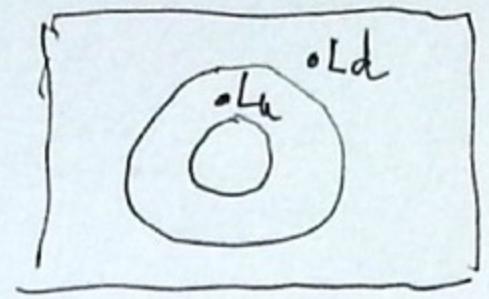
\rightarrow Thm 9.4 If L and \bar{L} are programmable, L is an algorithm.
 recursive



L, \bar{L} 의 정지 문제
 974 vs 624
 3x3 3!

Diagonalization Language L_d

TM is countable, Σ^* is countable
 (program) $\therefore (M_i, w_i)$ is also countable
 (input data)



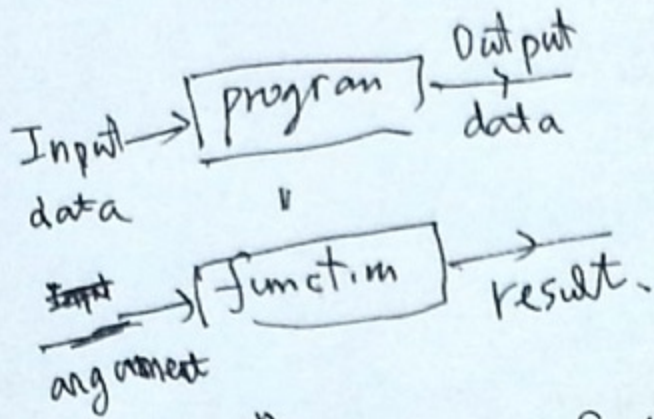
$$L_d = \{ w_i \in \Sigma^* \mid w_i \notin L(M_i) \}$$

$$L_u = \{ w_i \in \Sigma^* \mid w_i \in L(M_i) \}$$

$$L_d = \overline{L_u} \quad (\overline{L_u} = L_d)$$

L_d is NOT R.E.
 L_u is R.E.

Chap. 9-1 Computability (Partial Recursive Function)



$$f: \mathbb{N}^m \rightarrow \mathbb{N}^n \quad \text{vs} \quad f: \mathbb{N} \rightarrow \mathbb{N}$$

primitive recursive ftn.

zero ftn. $f: \mathbb{N}^m \rightarrow \{0\}$

$\vec{x} \in \mathbb{N}^m$
 (a_1, a_2, \dots, a_m)