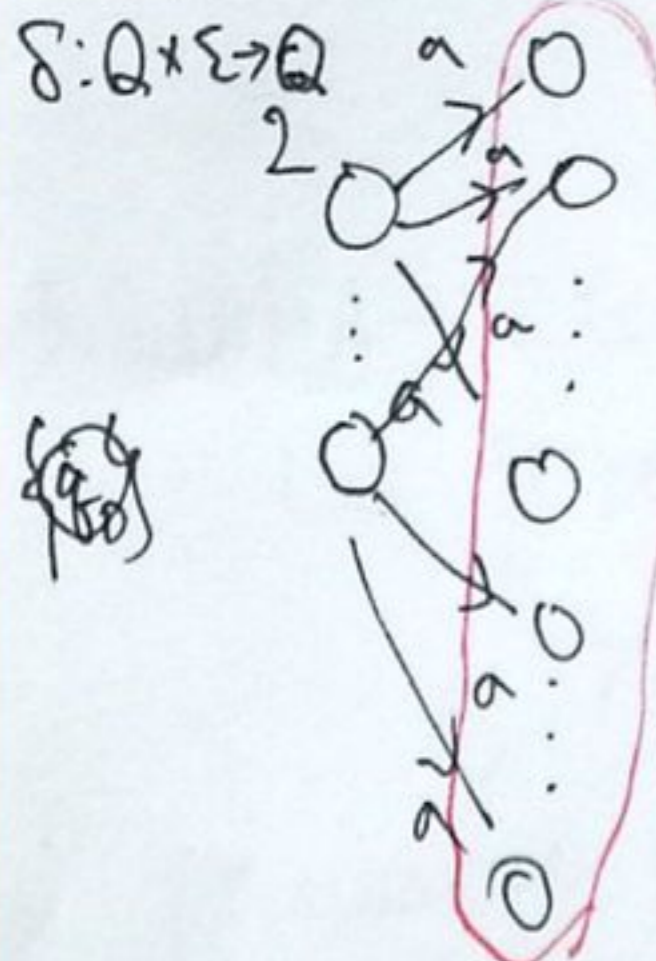
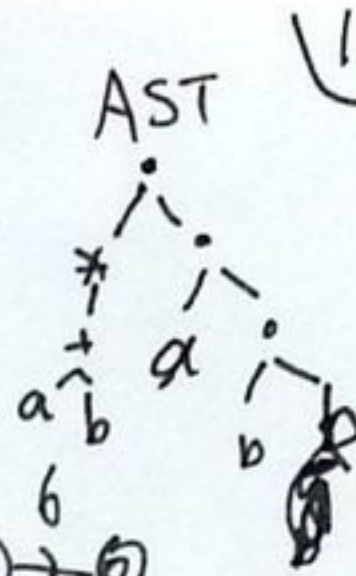
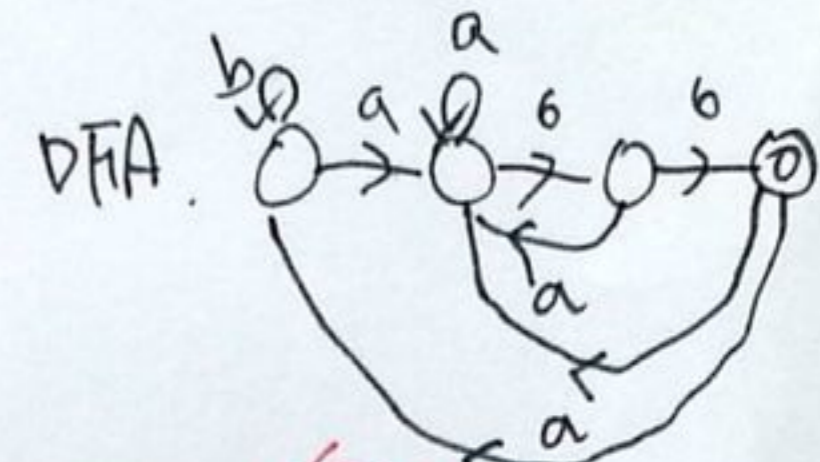
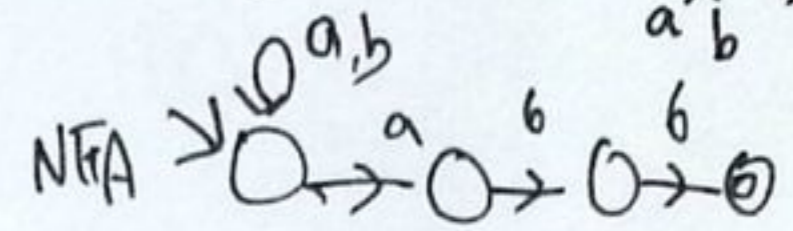


NFA  $\rightarrow$  DFA

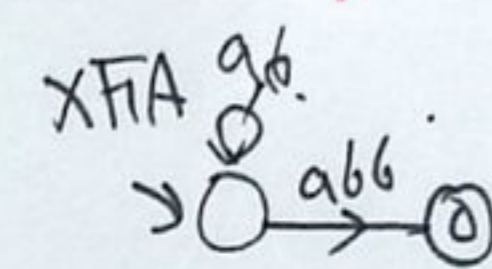
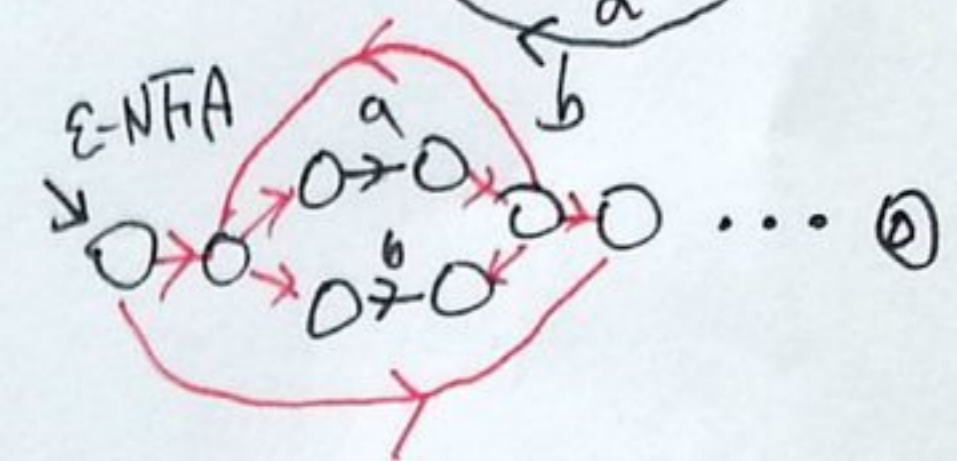
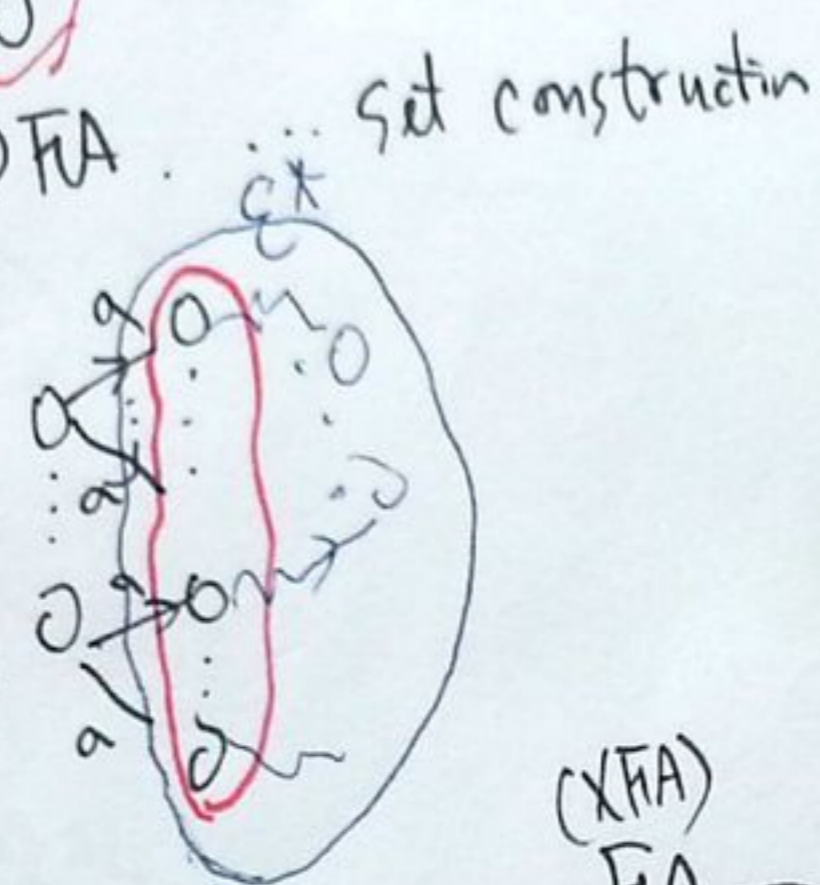
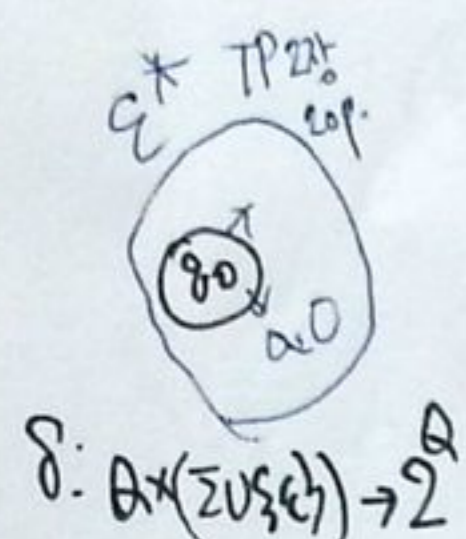


Set construction (XFA)  
 A set of states in NFA  
 is equ becomes  
 a state in DFA

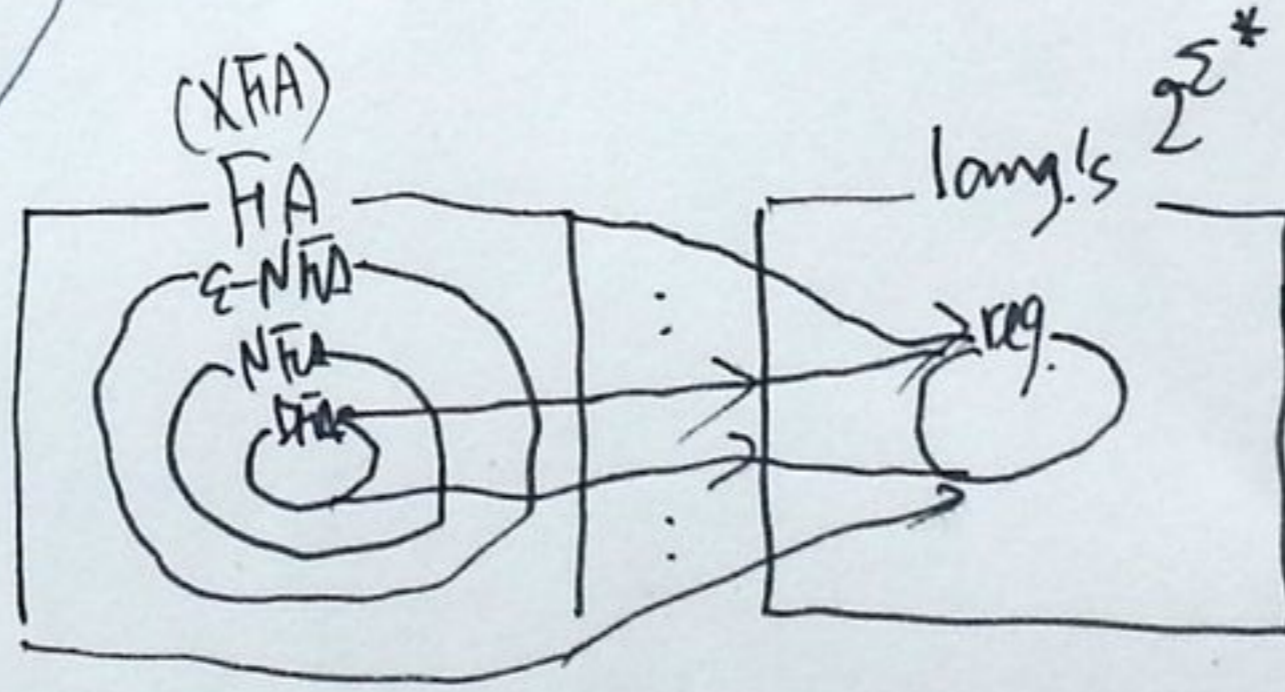
R.E.  $(a|b)^* abb$



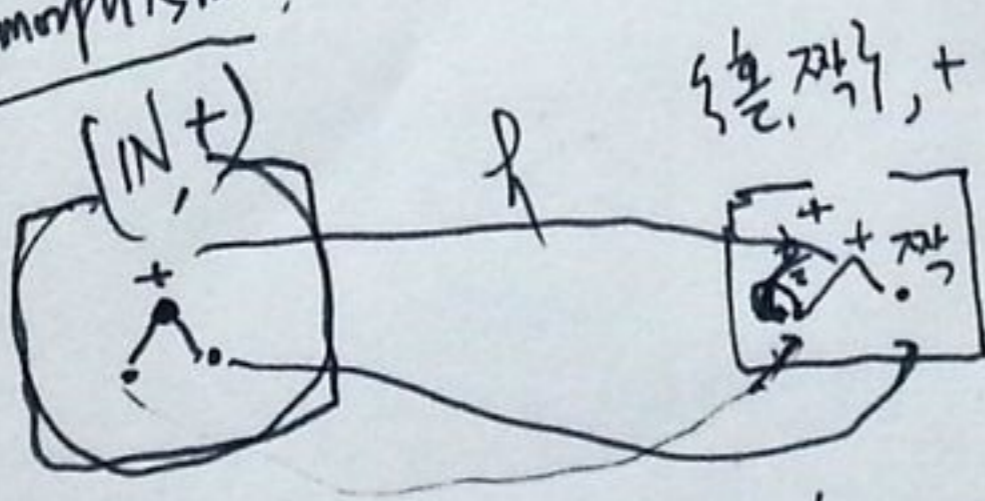
$\epsilon$ -NFA  $\rightarrow$  DFA



XNFA  
 $\delta: Q \times \Sigma^* \rightarrow 2^Q$



homomorphism:



총	합	정	2	정
정	합	정	정	정

abstract interpretation  
 concrete domain - abstract dom.

# Chap. 3 Regular expressions

syntax  $R \rightarrow a | \epsilon | \phi | R+R | R \cdot R | R^* | (R)$

semantics  $\{a\} \{ \epsilon \} \{ \phi \} L(R) \cup L(R') L(R) \cdot L(R') (L(R))^* L(R)$

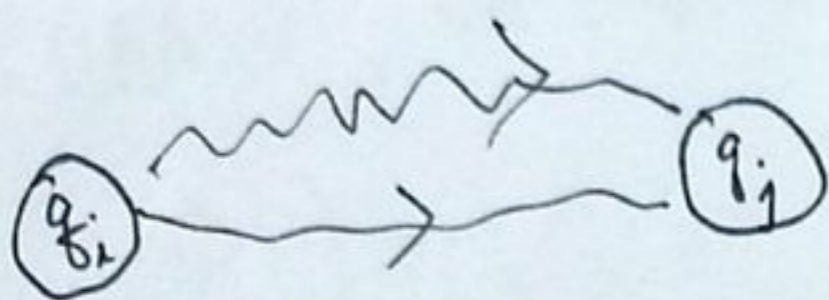
## 3.2. DFA $\rightarrow$ R.E

Let  $D$  be a

"  
DFA.

이 다르지 않?  
 $R \cdot \epsilon = \epsilon \cdot R = R$   
 $R \cdot \phi = \phi \cdot R = \phi$

1. recursive generation of R.E



$R_{ij}^0, R_{ij}^k$

2. 연결 방식