

15주-목 마지막 강의

- Cook's Theorem & Q & A

모든 complete(자)

NP 문제가 P 내에 해결되는가? - 어떤 NP 문제는 P에 해결된다.

- 모른다.
- 어떻게 하면 아나?

Cook의 정리에 따르면 NP-완성이 P에 해결되면 모든 NP 문제는 P에 해결된다.
- 하지만 NP-완성 문제를 해결하는 데 $O(n^2)$ 시간

Thm 10.9 Cook의 정리

n-SAT는 NP-complete이다.

Def. NP complete, P.

- 1. $P \in NP$
- 2. $\forall P' \in NP\text{-complete}, P' \leq_{PTR} P$.

Def. Satisfiability 문제 (SAT 문제)
(만족?)

Def. Boolean exp. (불대수식)

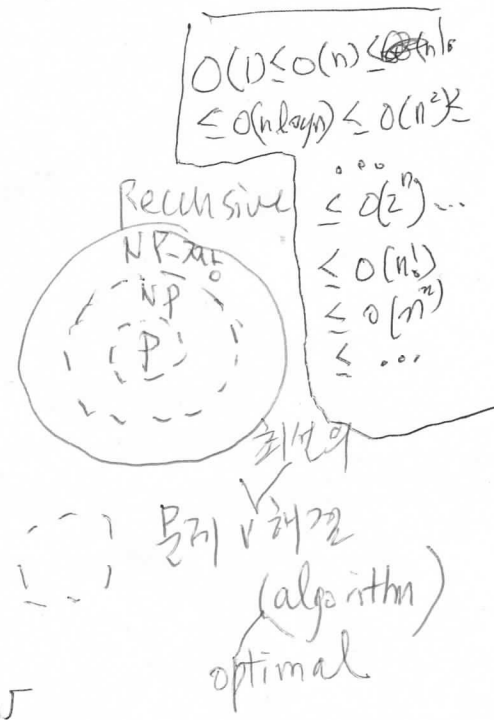
$b \rightarrow b \vee b \mid b \wedge b \mid \neg b \mid (b) \mid t \mid f$ 문제

Def. SAT with order n

$B(n) \dots$ 불대수식 B 의 변수 v_1, \dots, v_n 의 값으로 n 개의 리터럴

truth assignment of $B(v_1, \dots, v_n)$
 $B(v_1, \dots, v_n) = t$

$v_1 = t/f, v_2 = t/f, \dots, v_n = t/f$



이 truth assignment
= (문) 문제

$f(x, y)$
 $f(0, 1)$
등식: $f(0) = 0$
타입

(proof)

1) SAT ∈ NP.

1) $\forall P \in NP, \exists A P < n\text{-SAT}$

① truth-false assign t: 2^n
 각각의 해답이 있는지 P 시간.
 이것은 N time solution
 $n\text{-SAT} \in NP$

$(n \rightarrow B)$
 $\{1, \dots, n\} \rightarrow \{t, f\}$
 $\{b_1, \dots, b_n\}$

②

α_0 :	X_{00}	X_{01}	X_{02}		$X_{0, p(n)}$							
α_1 :	X_{10}	X_{11}	X_{12}		$X_{1, p(n)}$							
				<table border="1"> <tr> <td>$X_{i,j-1}$</td> <td>$X_{i,j}$</td> <td>$X_{i,j+1}$</td> </tr> <tr> <td>$X_{i+1,j-1}$</td> <td>$X_{i+1,j}$</td> <td>$X_{i+1,j+1}$</td> </tr> </table>	$X_{i,j-1}$	$X_{i,j}$	$X_{i,j+1}$	$X_{i+1,j-1}$	$X_{i+1,j}$	$X_{i+1,j+1}$		
$X_{i,j-1}$	$X_{i,j}$	$X_{i,j+1}$										
$X_{i+1,j-1}$	$X_{i+1,j}$	$X_{i+1,j+1}$										
$\alpha_{p(n)}$	$X_{p(n),0}$	$X_{p(n),1}$	$X_{p(n),2}$		$X_{p(n), p(n)}$							

$P(n) \triangleq$
 $p(\alpha_1, \alpha_2, \dots, \alpha_n)$
 P : polynomial
 (degree 3)

$(n \times (p(n) + 1))^2$ cells $\leq 10^6$ $3+3=6 \leq 10^6$

$X_{i,j}$	$X_{i,j+1}$	$X_{i,j+2}$
$X_{i+1,j}$	$X_{i+1,j+1}$	$X_{i+1,j+2}$

$\alpha_0 \Rightarrow \alpha_1 \Rightarrow \dots \Rightarrow \alpha_{p(n)}$
 $P(n)$ Time comp.
 $P(n)$ space comp.

\rightarrow N time
 M space

~~1)~~ $X_{i,j} \in \{0,1\}$ (Conf. TM $\otimes \Gamma^* \times \{0,1\}^*$)

$Y_{i,j} \in A \triangleq (X_{i,j} = A)$... 항상 177.

$$E_{m, v_0} = UASANA \bar{F}$$

$M \in TM's, w \in \Sigma^*$

$$(\exists y_{i,j}^A \vee \exists y_{i,j}^B)$$

$A \neq B$

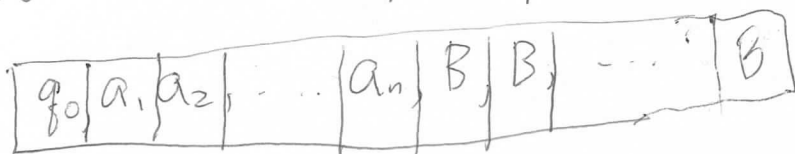
① Unique $0 \leq i, j \leq P(m)$

$$\parallel \Delta \forall i, j \in P(m)$$

$\bigwedge \neg (y_{i,j}^A \wedge y_{i,j}^B)$ when $A, B \in QUT$

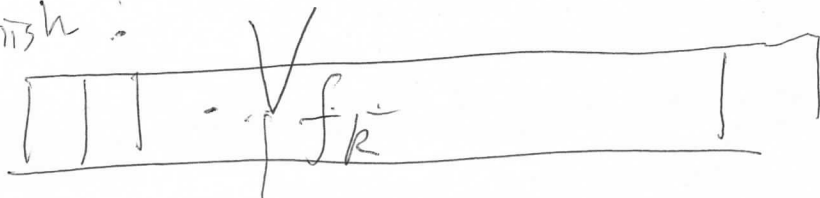
하나의 cell은 두개의 값으로 가질 수 없음

② Start! ... See TPR p 8, 9, 10



$$0 \leq k \leq P(m)$$

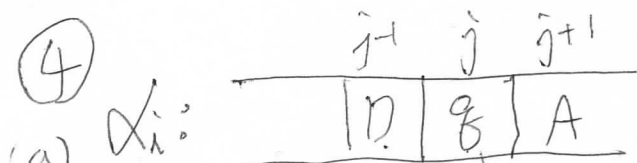
③ Finish!



$$f_k \in F$$

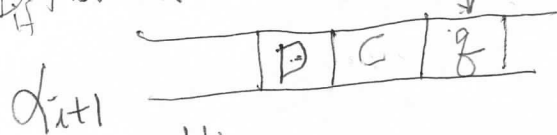
있음 \exists .

④

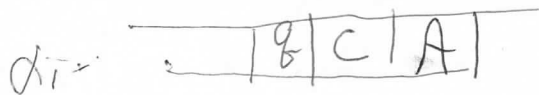


$$D, A \in P, q \in Q$$

① Move ~~left~~ right



② Move ~~right~~ left



(b) 다른 cell 에는 아무나 적시라!

⋮

결론 - SAT는 최초의 NP-강이다!