

A **Moore machine** is a six-tuple  $M_o = (Q, T, \Pi, \delta, \lambda, q_0)$  where

$Q$  is a set of states

$T$  is a set of input symbols

$\delta: Q \times T \rightarrow Q$

$\Pi$  is a set of **output** symbols,

$\lambda: Q \rightarrow \Pi$  is a **output** function from **state** to **output symbol**., and

$q_0 \in Q$  is a **start state**.

$Q, \Sigma, \delta$  and  $q_0$  are same as DFA but no final states( $F$ )

Assume input string is  $x = a_1 a_2 \dots a_n \in T^* (n \geq 0)$  and

$1 \leq \forall i \leq n: \delta(q_{i-1}, a_i) = q_i$ . Then

**Output string of Moore machine  $M_o$  for input string  $x$  is**

$\Lambda(M_o, x) = \lambda(q_0) \lambda(q_1) \dots \lambda(q_n) \in \Pi^*$ .

$|\Lambda(M_o, x)| = |x| + 1 = n + 1$ .

A **Mealy machine** is a six-tuple  $M_e = (Q, T, \Pi, \delta, \lambda, q_0)$  where

$Q$  is a set of states,

$T$  is a set of input symbols,

$\Pi$  is a set of **output** symbols,

$\delta: Q \times T \rightarrow Q$ ,

$\lambda: Q \times T \rightarrow \Pi$  is a **output** function

from **state transition** to **output** symbols, and

$q_0 \in Q$  is a **start** state.

Assume input string is  $x = a_1 a_2 \dots a_n \in T^*$  ( $n \geq 0$ ) and

$1 \leq \forall i \leq n: \delta(q_{i-1}, a_i) = q_i$ . Then

**Output string** of Mealy machine  $M_e$  for input string  $x = a_1 a_2 \dots a_n \in T^*$ .

$\Lambda(M_e, x) = \lambda(q_0, a_1) \lambda(q_1, a_2) \dots \lambda(q_{n-1}, a_n) \in \Pi^*$ .

$|\Lambda(M_e, x)| = |x| = n$ .

*What are the elements of the output vocabulary  $\Pi$ ?  
program segments or functions, ...*

*한글 모아쓰기 automata*

*FA is type 3 but Moore and Mealey machines are  
type 0(TM), if  $\Pi$  is a set of program segments.*