

12/5 (K) 2125 Turing - Church's Thesis.

Computability

Turing Machine - computer
 μ - recursive partial functions

Thm 4.2 Thm 4.3
 μ - primitive recursive functions
 + minimization (μ)

Chap 9. Undecidability

A language that is not R.E.

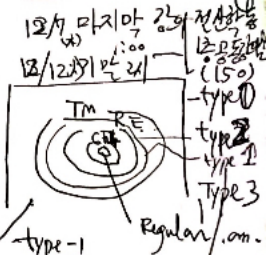
Halting Problem

Diagonalization Language L_d

Power set of Integer

Russell's paradox

Gödel's Incompleteness Theorem



recursive
 no loop forever
 always halt

TM ... countable M_i $i \in \mathbb{N}$
 numbering

Σ^* ... countable $w_i \in \Sigma^*$ ($i \in \mathbb{N}$)
 (M_i, w_i) pair

$L_d = \{w_i \in \Sigma^* \mid w_i \notin L(M_i)\}$

Assume $L_d = L(\underline{M})$. $\exists i \in \mathbb{N} \ M = M_i$
 $\exists i \in \mathbb{N} \ M = M_i$

$\exists i \in \mathbb{N} \ \therefore L_d \neq R.E. \ \text{of } M_i$

Self contradiction ... Paradox of Self Recursion!

non R.E.
 uncountable

Consider $L_u = \{w_i \in \Sigma^* \mid w_i \in L(M_i)\}$

$L_d = \overline{L_u}$