

# 3. Functions, Sequences, and Relations

## 3.1 Functions

A **function**(mapping, transformation)  $f$  from  $X$  to  $Y$  is  
an assignment of each(**all**) element of  $X$  **exactly one** element of  $Y$ .

**Def. 3.1.1** Let  $X$  and  $Y$  be sets. A function  $f \subseteq X \times Y$ (**relation**)  $\exists$ .

$$\forall x \in X, \exists \mathbf{1}y \in Y, \exists. (x, y) \in f.$$

$X$ : **domain**( 정의역 ) of the function  $f$

$Y$ : **codomain**( 공역 ) of the function  $f$

$\{y \in Y \mid \forall x \in X, (x, y) \in f\} \subseteq Y$  **range**( 치역 ) of the function  $f$

$f: X \rightarrow Y$   $f$  is a function from the **domain**  $X$  to **codomain**  $Y$

We write  $f(x) = y$  instead of  $(x, y) \in f$ .

A **function** is a **relation** ( 관계 ) satisfying two **conditions**

- 1) **total**: for all elements of  $X$  (domain)
- 2) **unique**: exactly one elements of  $Y$  (codomain)

**Def. 3.1.10** Let  $x \in \mathbf{Z}$  and  $n \in \mathbf{Z}^+$ . Then

$x \bmod n =$  positive remainder when  $x$  is divided by  $n$ .

$$\text{mod}: \mathbf{Z} \times \mathbf{Z}^+ \rightarrow \mathbf{Z}^+.$$

We may write  $\text{mod}(x, n) = m$  instead of  $x \bmod n = m$ .

**Def. 3.1.16** The **floor** and **ceiling** function:  $\lfloor \cdot \rfloor \lceil \cdot \rceil : \mathbf{R} \rightarrow \mathbf{Z}$ ,

$\lfloor x \rfloor = n \in \mathbf{Z}$ ,  $n$  is the **greatest integer** such that  $n \leq x$ . (**floor**)

$\lceil x \rceil = n \in \mathbf{Z}$ ,  $n$  is the **least integer** such that  $n \geq x$ . (**ceiling**)

Let  $x \in \mathbf{Z}$  and  $y \in \mathbf{Z}^+$ . Then

if  $x \geq 0 \rightarrow x \div y = \lfloor x/y \rfloor + x \bmod y$

/  $x \leq 0 \rightarrow x \div y = \lceil x/y \rceil + x \bmod y$

**fi**

**Def. 3.1.21** Let  $f: X \rightarrow Y$ .  $f$  is **one-to-one** (1-1 or **injective**; 단사 ), if

$$\forall x \in X \forall y \in Y: [(x \neq y) \Rightarrow (f(x) \neq f(y))], \text{ or } \textit{logically equivalent}$$

$$\forall x \in X \forall y \in Y: [(f(x) = f(y)) \Rightarrow (x = y)]. \quad \textit{contrapositive}(\text{ 대우 })$$

An **injective function** is called **injection**.

**Thm. 3.1.21'** Let  $f: X \rightarrow Y$ .  $f$  is **1-1**  $\Rightarrow |X| \leq |Y|$ .

**proof** Since  $f$  is **1-1**,  $|X| \stackrel{1-1}{=} |\text{range}(f)| \leq |\text{codom}(f)| = |Y| \therefore |X| \leq |Y|$ .

**Def. 3.1.28** 7 Let  $f: X \rightarrow Y$ .  $f$  is **onto** (or **surjective**; 전사 ), iff

$$\forall y \in Y \exists x \in X: (f(x) = y), \quad \text{or} \quad f(X) = Y (\text{range} = \text{codomain}).$$

A **surjective function** is called **surjection** (or **correspondence**).

**Thm. 3.1.28'** Let  $f: X \rightarrow Y$ .  $f$  is **onto**  $\Rightarrow |X| \geq |Y|$ .

**proof** Since  $f$  is **onto**,  $|X| \geq |\text{range}(f)| \stackrel{\text{onto}}{=} |\text{codom}(f)| = |Y| \therefore |X| \geq |Y|$ .

**Def. 3.1.34** Let  $f: X \rightarrow Y$ . If  $f$  is both **1-1** and **onto** then  $f$  is called **one-to-one onto** (or **bijective**; 전단사 ).

A **bijective function** is called **bijection** or **one-to-one onto function** or **one-to-one correspondence**; 짝짓기 ).

**Thm. 3.1.34'** Let  $f: X \rightarrow Y$  and  $f$  be **one-to-one** and **onto**(**bijective**).

The **inverse function** of  $f$ , denoted  $f^{-1}: Y \rightarrow X$ , is also is a **bijection**,

$f^{-1} = \{(y, x) \mid x \in X, f(x) = y \in Y\}$ . Then

$$\forall y \in Y: \exists! f^{-1}(y) = x \in X .\exists. f(x) = y.$$

**Thm. 3.1.34''** Let  $f: X \rightarrow Y$  and  $f$  be **one-to-one** and **onto**(**bijective**). Then

$$|X| = |Y|.$$

**Def. 3.1.40** Let  $g: A \rightarrow B$  and  $f: B \rightarrow C$ . The **composition** of  $f$  and  $g$ , denoted by  $f \circ g: A \rightarrow C$ , is defined by

$$(f \circ g)(a) = f(g(a)) \text{ or } f \circ g = \{(a, c) \mid f(a) = b, g(b) = c\}.$$

**Def. 3.1.40'** **Identity function(relation)** on  $A$

$$\iota_A: A \rightarrow A \text{ .}\exists. \iota_A = \{(a, a) \mid a \in A\} \text{ or } \forall a \in A: \iota_A(a) = a.$$

**Thm. 3.1.40''** Let  $f: A \rightarrow A$ . Then

$$f \circ \iota_A = \iota_A \circ f = f.$$

$\iota_A$  is a **identity** function for composition of functions on  $A$ .

**Def. 3.1.40'''** Let  $f_1, f_2: A \rightarrow \mathbf{R}$ .  $f_1 + f_2, f_1 f_2: A \rightarrow \mathbf{R}$  is defined by

$$(f_1 + f_2)(x) = f_1(x) + f_2(x) \text{ and } f_1 f_2(x) = f_1(x)f_2(x).$$

## Extension of Set Equivalence and Cardinality Revisited

We say that two sets  $A$  and  $B$  are **isomorphic** with respect to  $f$ ,  $A \cong_f B$ .

If  $f$  is a **bijection** from  $A$  to  $B$ ,  $f: A \leftrightarrow B$ . ( 짝짓기 )

$$\forall a \in A: \exists_1 f(a) \in B \text{ and } \forall b \in B \exists_1 f^{-1}(b) \in A.$$

We can identify  $B$  with  $A$  and  $f$ , and identify  $A$  with  $B$  and  $f^{-1}$  (vice versa)

**Set Isomorphism**

**Extended Set Equivalence**

**Def. Cardinality of Set, revisited**

Let  $A$  and  $B$  are sets. We say the **cardinalities** of  $A$  and  $B$  are same,

$|A| = |B|$ , if there is a **bijection**  $f: A \leftrightarrow B$ .

## Operation vs. Function

*Operation is a function on same sets.*

**Def. 3.1.45** Let  $f: X \times X \rightarrow X$ . Then  $f$  is called **binary operation** on  $X$ .

$\forall a, b, c \in X$ , we may write  $a f b = c$  instead of  $f(a, b) = c$ .  
 $f$  is the **infix binary operator**.

**Def. 3.1.48** Let  $f: X \rightarrow X$ . Then  $f$  is called **unary operation** on  $X$ .

$\forall a, b \in X$ , we may write  $f a = b$  instead of  $f(a) = b$ .  
 $f$  is the **prefix unary operator**.

**Def. 3.1.48'** Let  $f: X^n \rightarrow X$ . Then  $f$  is called *n-ary operation on X*.

$$\forall a_1, a_2, \dots, a_n, b \in X, f(a_1, a_2, \dots, a_n) = b \in X,$$

$f$  is the **prefix n-ary operator**.

**Def. 3.1.48''** Let  $f: X^n \rightarrow X^m$ . Then  $f$  is called *n-ary operation on X*.

$$\forall a_1, a_2, \dots, a_n \in X, f(a_1, a_2, \dots, a_n) = (b_1, b_2, \dots, b_m) \in X^m,$$

$f$  is called **multiple value returning function**

*Global variable considered harmful*

*Choe, Kwang-Moo*

*GOTO statement considered harmful*

*E. W. Dijkstra*



## 3.2 Sequences and Strings

**Sequence** a function whose domain is a set of numbers.

$a: \mathbf{N} \rightarrow A.$        $\mathbf{N}$  is a set of numbers       $A$  is a set.

We write  $a_n$  instead of  $a(n)$ .

$n$  is called **index** of the sequence  $a_n$

$\mathbf{N}$  is called **index set** of the sequence  $a_n$ .

$$\{a_n\} = \{a_n \mid n \in \mathbf{N}\} = \{a_n\}_{n \in \mathbf{N}}.$$

### Some Useful Sequences

*polynomial sequences*       $n^2, n^3, n^4, \dots, n^k, \dots$

*exponential sequences*       $2^n, 3^n, \dots, k^n, \dots, n!, \dots, n^n, \dots$

**Definition 6** Let  $\{s_n\}_{n \in \mathbf{N}}$  ( $s: \mathbf{N} \rightarrow A$ ) be a sequence over  $A$  and  $\leq$  is a **partial order** on  $A$  (See 3.3). Then

If  $s_n < s_{n+1}$ ,  $s$  is **increasing** or (**strictly increasing**),  
 If  $s_n > s_{n+1}$ ,  $s$  is **decreasing** or (**strictly decreasing**),  
 If  $s_n \leq s_{n+1}$ ,  $s$  is **nondecreasing** or (**increasing**), and  
 If  $s_n \geq s_{n+1}$ ,  $s$  is **nonincreasing** or (**decreasing**).

**Def. 3.2.11** Let  $\{s_n\}_{n \in \mathbf{N}}$  be a sequence.

If  $\mathbf{I} \subseteq \mathbf{N}$ , then  $\{s_i\}_{i \in \mathbf{I}}$  is the **subsequence** of  $\{s_n\}_{n \in \mathbf{N}}$ .

**Def. 3.2.14**

$$\sum_{i=m}^n a_i = a_m + a_{m+1} + \dots + a_n = \sum_{i \in \mathbf{N}_{m,n}} a_i, \quad \text{sum}(\text{sigma}) \text{ notation}$$

$$\prod_{i=m}^n a_i = a_m \cdot a_{m+1} \cdot \dots \cdot a_n = \prod_{i \in \mathbf{N}_{m,n}} a_i, \quad \text{product}(\text{pi}) \text{ notation}$$

where  $\mathbf{N}_{m,n} = \{m, m+1, \dots, n\}$  ( $m \leq n$ ).

$m$ : lower limit     $n$ : upper limit     $i$ : index variable

$\mathbf{N}_{m,n}$ : index set

**Def. 3.2.10** A **String** over  $V$ .

Let  $n \geq 0$ ,  $s: \{1, 2, \dots, n\} \rightarrow V = \{a, b, \dots, z\}$ .

We write  $s = (b, o, y)$  or  $s = \text{boy}$  for short

instead of  $s(1) = b, s(2) = o, s(3) = y$ .

$s$  is called the finite **string** over  $V$  of length  $n$ .

$V$  is called the **vocabulary(alphabet)** of string  $s$ .

**length** of a string  $|\alpha|$   
 number of elements in  $\alpha$ .

**null(empty)** string  $\lambda$  or  $\varepsilon$

**Recursive definition of  $V^n$ .**

**basis**  $V^0 = \{\lambda\}$

**recursion**  $V^n = V^{n-1} \cdot V$

**universe of string over  $V$**

$$V^* = V^0 \cup V^1 \cup V^2 \cup \dots$$

$$\text{where } V^0 = \{\lambda\}, V^1 = V, V^2 = V \cdot V, \dots, V^n = V^{n-1} \cdot V, \dots$$

$$\forall \alpha \in V^*, \alpha \cdot \varepsilon = \varepsilon \cdot \alpha = \alpha.$$

$\varepsilon$  is a **identity element on concatenation**.

**concatenation of two string  $\alpha$  and  $\beta$  ( $\alpha, \beta \in X^*$ )**

$$\therefore X^* \times X^* \rightarrow X^*$$

$$\alpha \cdot \beta = \alpha\beta$$

$$|\alpha \cdot \beta| = |\alpha| + |\beta|$$

$$\text{school} \cdot \text{boy} = \text{schoolboy}$$

**Def 3.2.26** Let  $\alpha = \gamma\beta\delta$  where  $\alpha, \beta, \gamma, \delta \in V^*$ .

Then  $\beta$  is a **substring** of  $\alpha$ ,  $\gamma$  is a **prefix** of  $\alpha$ , and  $\delta$  is a **postfix** of  $\alpha$ .

Ex.  $\alpha = \text{boy}$ .  $\text{substring}(\alpha) = \{\lambda, b, o, y, bo, oy, boy\}$

$\text{prefix}(\alpha) = \{\lambda, b, bo, boy\}$

$\text{posfix}(\alpha) = \{\lambda, y, oy, boy\}$

## 한글

A set of **strings** over  $\Sigma_{24} = \{ \neg, \perp, \dots, \bar{\circ} \} \cup \{ \vdash, \vDash, \dots, \lceil \}$ .

My MS dissertation  $\Sigma_{29} = \Sigma_{24} \cup \{ \neg\neg, \vDash\vDash, \text{ㅁㅁ}, \text{ㅂㅂ}, \text{ㅅㅅ} \}$ .

computer keyboard  $\Sigma_{33} = \Sigma_{29} \cup \{ \text{ㅈ}, \text{ㅊ}, \text{ㅋ}, \text{ㆁ} \}$ .

삼성 hp  $\Sigma_{11} = \{ \lceil, \cdot, \text{—}, \neg, \perp, \vDash, \text{ㅁ}, \text{ㅂ}, \text{ㅅ}, \text{ㅈ}, \text{ㅊ}, \text{ㅇ}, \text{blank} \}$ .

LG hp  $\Sigma_{12} = \{ \vdash, \perp, \text{—}, \lceil, \neg, \perp, \vDash, \text{ㄱ}, \text{ㄴ}, \text{ㅅ}, *, + \}$ .

한글음절 (unicode) =  $19 \times 21 \times 28 = 1,1172$

### 3.3 Relations

**Def. 3.3.1** If  $R \subseteq X \times Y$ , then  $R$  is called a **(binary) relation** from  $X$  to  $Y$ .

We may write  $a R b$ , if  $(a, b) \in R$ .

Let  $R \subseteq X \times X$ ,  $R$  is called a **relation on  $X$** .

A **digraph**(directed graph)  $G = (V, E)$  where  $E \subseteq V \times V$ .

- 1)  $V$ : a set of vertices
- 2)  $E \subseteq V \times V$ : a set of edges

Let  $R$  be a relation on  $X$  ( $R \subseteq X \times X$ ). Then

**Def. 3.3.6**  $R$  is **reflexive**, if  $\forall x \in X: (x, x) \in R$  (or  $x R x$ ). or  $\iota_X \subseteq R$ .

**Def. 3.3.6'**  $R$  is **irreflexive**, if  $\forall x \in X: (x, x) \notin R$  (or  $x \not R x$ ). or  $R \cap \iota_X = \emptyset$ .

**Def. 3.3.9**  $R$  is **symmetric**, if  $\forall x, y \in X: (x, y) \in R \Rightarrow (y, x) \in R$ .

(or  $x R y \Rightarrow y R x$ ) or  $R = R^{-1}$ .

**Def. 3.3.12**  $R$  is *antisymmetric*, if  $\forall x, y \in X: (x, y) \in R \wedge (y, x) \in R \Rightarrow x=y$ .  
 (or  $((x R y) \wedge (x \neq y)) \Rightarrow (y \not R x)$  or  $R \cap R^{-1} \subseteq \iota_X$ .)

**Thm. 3.3.12'** If  $R$  is *antisymmetric*,  
 then  $R$  is not (always) *reflexive* nor *irreflexive*.

**Def. 3.3.12''**  $R$  is *asymmetric*, if  $\forall x, y \in X, (x, y) \in R \Rightarrow (y, x) \notin R$ .  
 (or  $x R y \Rightarrow y \not R x$  or  $R \cap R^{-1} = \emptyset$ .)

**Thm. 3.3.12'''** If  $R$  is *asymmetric*,  
 then  $R$  is *irreflexive* and *antisymmetric*.

**Def. 3.3.13**  $R$  is *transitive*, if  $\forall x, y, z \in X: (x, y) \in R \wedge (y, z) \in R \Rightarrow (x, z) \in R$ .  
 (or  $x R y \wedge y R z \Rightarrow x R z$  or  $R \cdot R \subseteq R$ .)



**Def. 3.3.20** Let  $R \subseteq X \times X$ . Then  $R$  is a **partial order**,  
if  $R$  is **reflexive**, **antisymmetric**, and **transitive**.

**Def. 3.3.20'** Let  $\leq \subseteq X \times X$ . Then  $\leq$  is a **partial order**. Then  
 $(X, \leq)$  is called a **poset** (or **partially ordered sets**).

If  $(x \leq y) \vee (y \leq x)$ , then  $x$  and  $y$  are **comparable**.

Otherwise  $((x \not\leq y) \wedge (y \not\leq x))$ , **incomparable**.

If  $\forall x, y \in X$ ,  $x$  and  $y$  are comparable  $\leq$  is **total** (or **linear**) order.

**Exa. 3.3.21** Let  $R \subseteq \mathbf{Z}^+ \times \mathbf{Z}^+$ . and  $(x, y) \in \mathbf{Divides}$ , if  $x$  divides  $y$ . Then  
 $(\mathbf{Z}^+, \mathbf{Divides})$  is a **poset**.

$(\mathbf{Z}^+, \leq)$  is a **total order**.

$(2^S, \subseteq)$  is a **poset**.

## Hasse Diagram

Let  $(A, \leq)$  be a poset. We say  $a$  **covers**  $b$ , if  $a < b \wedge \neg \exists c \in A, a \leq c \wedge c \leq b$ .  
 $(A, \text{covers})$  is a Hasse Diagram.

## Maximal and Minimal Elements

Let  $(A, \leq)$  be a poset.

We say  $a$  is a **maximal**, if  $\neg \exists b \in A, a < b$ .

We say  $a$  is a **minimal**, if  $\neg \exists b \in A, b < a$ .

We say  $a$  is a **greatest element**, if  $\forall b \in A, b \leq a$ .

We say  $a$  is a **least elements**, if  $\forall b \in A, a \leq b$ .

Let  $S \subseteq A$ .  $u \in A$  is called **upper bound** of  $S$ , if  $\forall a \in S, a \leq u$ .

$l \in A$  is called **lower bound** of  $S$ , if  $\forall a \in S, l \leq a$ .

$x \in A$  is called the **least upper bound** of  $S$ , if  $\forall a \in S, a \leq x$ ,

$x \leq z$ , if  $\forall z \in \text{upper bound of } S$ .

$y \in A$  is called the **greatest lower bound** of  $S$ , if  $\forall a \in S, y \leq a$ ,

$z \leq y$ , if  $\forall z \in \text{lower bound of } S$ .

## Lattice

A poset  $(A, \leq)$  is called **lattice**. If every pair of elements has both a **least upper bound**, and a **greatest lower bound**.

$(\mathbf{Z}, \leq)$  is a lattice.

$\min, \max: \mathbf{Z} \times \mathbf{Z} \rightarrow \mathbf{Z}$

$(\mathbf{Z}^+, |)$  is a lattice.

$\gcd, \text{lcm}: \mathbf{Z}^+ \times \mathbf{Z}^+ \rightarrow \mathbf{Z}^+$

$(2^S, \subseteq)$  is a lattice.

$\cap, \cup: 2^S \times 2^S \rightarrow 2^S$

**Def. 3.3.23** Let  $R \subseteq X \times Y$ , The **inverse** of  $R$ , denoted  $R^{-1} \subseteq Y \times X$ ,

$$R^{-1} = \{(y, x) \in Y \times X \mid (x, y) \in R\}$$

**Def. 3.3.25** Let  $R_1 \subseteq X \times Y$ , and  $R_2 \subseteq Y \times Z$ . The **composition** of  $R_1$  and

$R_2$ , denoted  $R_2 \circ R_1 \subseteq X \times Z$ ,

$$R_2 \circ R_1 = \{(x, z) \in X \times Z \mid (x, y) \in R_1, (y, z) \in R_2\}$$

**Three faces of the relation  $R$  from  $A$  to  $B$ :  $R \subseteq A \times B$ .**

**i)  $R$  is a set of pairs (subset of Cartesian product)**

$$R \subseteq A \times B, (a, b) \in R \text{ where } a \in A \text{ and } b \in B.$$

**ii)  $R$  is a infix binary boolean (relational) function**

$$R: A \times B \rightarrow \{\mathbf{T}, \mathbf{F}\}, a R b \text{ where } a \in A \text{ and } b \in B.$$

$R$  is rectangular boolean matrix. (See Sec. 3.5)

**iii)  $R$  is a function from  $A$  to  $2^B$ .**

$$R: A \rightarrow 2^B.$$

$$R(a) = \{b_1, b_2, \dots, b_n\} \text{ where } a \in A \text{ and } \{b_1, b_2, \dots, b_n\} \in 2^B.$$

$$\text{if } 1 \leq \forall i \leq n, (a, b_i) \in R \text{ or } a R b_i \text{ for } n \geq 0,$$

$$\text{if } n=0, R(a) = \{b_1, b_2, \dots, b_n\} = \emptyset.$$

Note that  $\forall a \in A, \exists \{b_1, b_2, \dots, b_n\} \in 2^B$ . (unique)

**set valued function.**

### 3.4 Equivalence Relations

**Thm 3.4.1** Let  $S$  be a partition of a set  $X$ . We define a relation on  $X$  as

$$R = \{(x, y) \in X \times X \mid x, y \in s, s \in S\}.$$

Then  $R$  is **reflexive, symmetric, and transitive**.

**proof**  $\forall x, y, z \in X: \exists s \in S. x, y, z \in s.$

$\therefore (x, x) \in R. \quad \therefore R$  is **reflexive**.

$\therefore (x, y) \in R \Rightarrow (y, x) \in R. \quad \therefore R$  is **symmetric**.

$\therefore (x, y), (y, x) \in R \Rightarrow (x, z) \in R. \quad \therefore R$  is **transitive**.

**Thm 3.4.3** A relation that is **reflexive, symmetric, and transitive** is called an **equivalent relation**.

**Exa.**  $= (\equiv \iota_{\mathbf{N}}) \subseteq \mathbf{N} \times \mathbf{N}$  is an **equivalent relation**

$< \subseteq \mathbf{N} \times \mathbf{N}$  is a **irreflexive total order**

$\leq (\equiv < \cup =) \subseteq \mathbf{N} \times \mathbf{N}$  is a **reflexive total order**

**Thm 3.4.8** Let  $R \subseteq X \times X$  be an **equivalent** relation on  $X$ , let

$[a]_R = \{x \in X \mid a R x\}$  **equivalent class** of  $X$  given by the relation  $R$ .

Then  $S = \{[a]_R \mid a \in X\}$  is a **partition** of  $X$ .

**proof** Since  $R$  is **reflexive**,  $\forall a \in X: a R a. \therefore a \in [a]_R$ .

If  $a R b$ , then  $[a]_R \subseteq [b]_R$ , since  $R$  is **transitive** and

$[b]_R \subseteq [a]_R$ , since  $R$  is **symmetric** and **transitive**.  $\therefore$  If  $a R b$   $[a]_R = [b]_R$ .

If  $a \not R b$ ,  $[a]_R \cap [b]_R = \emptyset$ .

<b>Relation</b>	$R \subseteq X \times X$	$ R  \leq n^2$	$O(n^2)$ where $ X  = n$ .
<b>Equivalent relation</b>	$S = \{[a]_R \mid a \in X\}$	$ S  \leq n$	$O(n)$ .

***3.5 Matrices of Relations***

***3.6 Relational Database***