

3. Functions, Sequences, and Relations

3.1 Functions

A **function**(mapping, transformation) f from X to Y is
an assignment of each(**all**) element of X **exactly one** element of Y .

Def. 3.1.1 Let X and Y be sets. A function $f \subseteq X \times Y$ (**relation**) \exists .

$$\forall x \in X, \exists! y \in Y, \exists (x, y) \in f.$$

X : domain(정의역) of the function f

Y : codomain(공역) of the function f

$\{y \in Y \mid \exists x \in X, (x, y) \in f\} \subseteq Y$ range(치역) of the function f

$f: X \rightarrow Y$ f is a function from the **domain** X to **codomain** Y

We write $f(x) = y$ instead of $(x, y) \in f$.

A **function** is a **relation** (관계) satisfying two **conditions**

- 1) **total**: for all elements of X (domain)
- 2) **unique**: exactly one elements of Y (codomain)

Def. 3.1.10 Let $x \in \mathbf{Z}$ and $n \in \mathbf{Z}^+$. Then

$x \bmod n =$ positive remainder when x is divided by n .

$$\text{mod}: \mathbf{Z} \times \mathbf{Z}^+ \rightarrow \mathbf{Z}^+.$$

We may write $\text{mod}(x, n) = m$ instead of $x \bmod n = m$.

Def. 3.1.16 The **floor** and **ceiling** function: $\lfloor \cdot \rfloor \lceil \cdot \rceil : \mathbf{R} \rightarrow \mathbf{Z}$,

$\lfloor x \rfloor = n \in \mathbf{Z}$, n is the **greatest integer** such that $n \leq x$. (**floor**)

$\lceil x \rceil = n \in \mathbf{Z}$, n is the **least integer** such that $n \geq x$. (**ceiling**)

Let $x \in \mathbf{Z}$ and $y \in \mathbf{Z}^+$. Then

$$\text{if } x \geq 0 \rightarrow x \div y = \lfloor x/y \rfloor + x \bmod y$$

$$/ x \leq 0 \rightarrow x \div y = \lceil x/y \rceil + x \bmod y$$

fi

Def. 3.1.21 Let $f: X \rightarrow Y$. f is **one-to-one** (1-1 or **injective**; 단사), if

$$\forall x \in X \forall y \in Y: [(x \neq y) \Rightarrow (f(x) \neq f(y))], \text{ or } \textit{logically equivalent}$$

$$\forall x \in X \forall y \in Y: [(f(x) = f(y)) \Rightarrow (x = y)]. \quad \textit{contrapositive}(\text{ 대우 })$$

An **injective function** is called **injection**.

Thm. 3.1.21' Let $f: X \rightarrow Y$. f is **1-1** $\Rightarrow |X| \leq |Y|$.

proof Since f is **1-1**, $|X| = |\text{codom}(f)| \leq |\text{range}(f)| = |Y| \therefore |X| \leq |Y|$.

Def. 3.1.28 7 Let $f: X \rightarrow Y$. f is **onto** (or **surjective**; 전사), iff

$$\forall y \in Y \exists x \in X: (f(x) = y), \quad \text{or} \quad f(X) = Y(\text{range} = \text{codomain}).$$

A **surjective function** is called **surjection** (or **correspondence**).

Thm. 3.1.28' Let $f: X \rightarrow Y$. f is **onto** $\Rightarrow |X| \geq |Y|$.

proof Since f is **onto**, $|X| \geq |\text{codom}(f)| = |\text{range}(f)| = |Y| \therefore |X| \geq |Y|$.

Def. 3.1.34 Let $f: X \rightarrow Y$. If f is both **1-1** and **onto** then f is called **one-to-one onto** (or **bijective**; 전단사).

A **bijective function** is called **bijection** or **one-to-one onto function** or **one-to-one correspondence**; 짝짓기).

Thm. 3.1.34' Let $f: X \rightarrow Y$ and f be **one-to-one** and **onto**(**bijective**).

The **inverse function** of f , denoted $f^{-1}: Y \rightarrow X$, is also is a **bijection**,

$f^{-1} = \{(y, x) \mid x \in X, f(x) = y \in Y\}$. Then

$$\forall y \in Y: \exists! f^{-1}(y) = x \in X .\exists. f(x) = y.$$

Thm. 3.1.34'' Let $f: X \rightarrow Y$ and f be **one-to-one** and **onto**(**bijective**). Then

$$|X| = |Y|.$$

Def. 3.1.40 Let $g: A \rightarrow B$ and $f: B \rightarrow C$. The **composition** of f and g , denoted by $f \circ g: A \rightarrow C$, is defined by

$$(f \circ g)(a) = f(g(a)) \text{ or } f \circ g = \{(a, c) \mid f(a) = b, g(b) = c\}.$$

Def. 3.1.40' **Identity function(relation)** on A

$$\iota_A: A \rightarrow A \text{ .}\exists. \iota_A = \{(a, a) \mid a \in A\} \text{ or } \forall a \in A: \iota_A(a) = a.$$

Thm. 3.1.40'' Let $f: A \rightarrow A$. Then

$$f \circ \iota_A = \iota_A \circ f = f.$$

ι_A is a **identity** function for composition of functions on A .

Def. 3.1.40''' Let $f_1, f_2: A \rightarrow \mathbf{R}$. $f_1 + f_2, f_1 f_2: A \rightarrow \mathbf{R}$ is defined by

$$(f_1 + f_2)(x) = f_1(x) + f_2(x) \text{ and } f_1 f_2(x) = f_1(x)f_2(x).$$

Extension of Set Equivalence and Cardinality Revisited

We say that two sets A and B are **isomorphic** with respect to f , $A \cong_f B$.

If f is a **bijection** from A to B , $f: A \leftrightarrow B$. (짝짓기)

$$\forall a \in A: \exists_1 f(a) \in B \text{ and } \forall b \in B \exists_1 f^{-1}(b) \in A.$$

We can identify B with A and f , and identify A with B and f^{-1} (vice versa)

Set Isomorphism

Extended Set Equivalence

Def. Cardinality of Set, revisited

Let A and B are sets. We say the **cardinalities** of A and B are same,

$|A| = |B|$, if there is a **bijection** $f: A \leftrightarrow B$.

Operation vs. Function

Operation is a function on same sets.

Def. 3.1.45 Let $f: X \times X \rightarrow X$. Then f is called **binary operation** on X .

$\forall a, b, c \in X$, we may write $a f b = c$ instead of $f(a, b) = c$.

f is the **infix binary operator**.

Def. 3.1.48 Let $f: X \rightarrow X$. Then f is called **unary operation** on X .

$\forall a, b \in X$, we may write $f a = b$ instead of $f(a) = b$.

f is the **prefix unary operator**.

Def. 3.1.48' Let $f: X^n \rightarrow X$. Then f is called *n-ary operation on X*.

$$\forall a_1, a_2, \dots, a_n, b \in X, f(a_1, a_2, \dots, a_n) = b \in X,$$

f is the **prefix n-ary operator**.

Def. 3.1.48'' Let $f: X^n \rightarrow X^m$. Then f is called *n-ary operation on X*.

$$\forall a_1, a_2, \dots, a_n \in X, f(a_1, a_2, \dots, a_n) = (b_1, b_2, \dots, b_m) \in X^m,$$

f is called **multiple value returning function**

Global variable considered harmful

Choe, Kwang-Moo

GOTO statement considered harmful

E. W. Dijkstra

3.2 Sequences and Strings

Sequence a function whose domain is a set of numbers.

$a: \mathbf{N} \rightarrow A.$ \mathbf{N} is a set of numbers A is a set.

We write a_n instead of $a(n)$.

n is called **index** of the sequence a_n

\mathbf{N} is called **index set** of the sequence a_n .

$$\{a_n\} = \{a_n \mid n \in \mathbf{N}\} = \{a_n\}_{n \in \mathbf{N}}.$$

Some Useful Sequences

polynomial sequences $n^2, n^3, n^4, \dots, n^k, \dots$

exponential sequences $2^n, 3^n, \dots, k^n, \dots, n!, \dots, n^n, \dots$

Definition 6 Let $\{s_n\}_{n \in \mathbf{N}}$ ($s: \mathbf{N} \rightarrow A$) be a sequence over A and \leq is a **partial order** on A (See 3.3). Then

If $s_n < s_{n+1}$, s is **increasing** or (**strictly increasing**),
 If $s_n > s_{n+1}$, s is **decreasing** or (**strictly decreasing**),
 If $s_n \leq s_{n+1}$, s is **nondecreasing** or (**increasing**), and
 If $s_n \geq s_{n+1}$, s is **nonincreasing** or (**decreasing**).

Def. 3.2.11 Let $\{s_n\}_{n \in \mathbf{N}}$ be a sequence.

If $\mathbf{I} \subseteq \mathbf{N}$, then $\{s_i\}_{i \in \mathbf{I}}$ is the **subsequence** of $\{s_n\}_{n \in \mathbf{N}}$.

Def. 3.2.14

$$\sum_{i=m}^n a_i = a_m + a_{m+1} + \dots + a_n = \sum_{i \in \mathbf{N}_{m,n}} a_i, \quad \text{sum}(\text{sigma}) \text{ notation}$$

$$\prod_{i=m}^n a_i = a_m \cdot a_{m+1} \cdot \dots \cdot a_n = \prod_{i \in \mathbf{N}_{m,n}} a_i, \quad \text{product}(\text{pi}) \text{ notation}$$

where $\mathbf{N}_{m,n} = \{m, m+1, \dots, n\}$ ($m \leq n$).

m : lower limit n : upper limit i : index variable

$\mathbf{N}_{m,n}$: index set

Def. 3.2.10 A **String** over V .

Let $n \geq 0$, $s: \{1, 2, \dots, n\} \rightarrow V = \{a, b, \dots, z\}$.

We write $s = (b, o, y)$ or $s = boy$ for short

instead of $s(1) = b$, $s(2) = o$, $s(3) = y$.

s is called the finite **string** over V of length n .

V is called the **vocabulary(alphabet)** of string s .

length of a string α $|\alpha|$
 number of elements in α .

null(empty) string λ or ε

concatenation of two string α and β ($\alpha, \beta \in X^*$)

$$\therefore X^* \times X^* \rightarrow X^*$$

$$\alpha \cdot \beta = \alpha\beta$$

$$|\alpha \cdot \beta| = |\alpha| + |\beta|$$

$school \cdot boy = schoolboy$

Recursive definition of V^n .

basis $V^0 = \{\lambda\}$

recursion $V^n = V^{n-1} \cdot V.$

universe of string over V

$$V^* = V^0 \cup V^1 \cup V^2 \cup \dots$$

where $V^0 = \{\lambda\}$, $V^1 = V$, $V^2 = V \cdot V$, ..., $V^n = V^{n-1} \cdot V$, , ...

$$\forall \alpha \in V^*, \alpha \cdot \varepsilon = \varepsilon \cdot \alpha = \alpha.$$

ε is a **identity element on concatenation**.

Def 3.2.26 Let $\alpha = \gamma\beta\delta$ where $\alpha, \beta, \gamma, \delta \in V^*$.

Then β is a **substring** of α , γ is a **prefix** of α , and δ is a **postfix** of α .

Ex. $\alpha = \text{boy}$. $\text{substring}(\alpha) = \{\lambda, b, o, y, bo, oy, boy\}$

$prefix(\alpha) = \{\lambda, b, bo, boy\}$

$posfix(\alpha) = \{\lambda, y, oy, boy\}$

한글

A set of **strings** over $\Sigma_{24} = \{ \neg, \perp, \dots, \bar{\circ} \} \cup \{ \vdash, \vDash, \dots, \lrcorner \}$.

My MS dissertation $\Sigma_{29} = \Sigma_{24} \cup \{ \pi, \sqcup, \boxplus, \bowtie, \boxtimes \}$.

computer keyboard $\Sigma_{33} = \Sigma_{29} \cup \{ \mathbb{H}, \mathbb{H}, \mathbb{K}, \mathbb{K} \}$.

삼성 hp $\Sigma_{11} = \{ |, \cdot, -, \neg, \perp, \sqcup, \boxplus, \boxtimes, \circ, \}$.

LG hp $\Sigma_{12} = \{ \vdash, \perp, -, \lrcorner, \neg, \perp, \sqcup, \boxplus, \boxtimes, \circ \}$.

3.3 Relations

Def. 3.3.1 If $R \subseteq X \times Y$, then R is called a **(binary) relation** from X to Y .

We may write $a R b$, if $(a, b) \in R$.

Let $R \subseteq X \times X$, R is called a **relation on X** .

A **digraph**(directed graph) $G = (V, E)$ where $E \subseteq V \times V$.

1) V : a set of vertices

2) $E \subseteq V \times V$: a set of edges

Let R be a relation on X ($R \subseteq X \times X$). Then

Def. 3.3.6 R is **reflexive**, if $\forall x \in X: (x, x) \in R$ (or $x R x$). or $\iota_X \subseteq R$.

Def. 3.3.6' R is **irreflexive**, if $\forall x \in X: (x, x) \notin R$ (or $x \not R x$). or $R \cap \iota_X = \emptyset$.

Def. 3.3.9 R is **symmetric**, if $\forall x, y \in X: (x, y) \in R \Rightarrow (y, x) \in R$.

(or $x R y \Rightarrow y R x$) or $R = R^{-1}$.

Def. 3.3.12 R is *antisymmetric*, if $\forall x, y \in X: (x, y) \in R \wedge (y, x) \in R \Rightarrow x=y$.
 (or $((x R y) \wedge (x \neq y)) \Rightarrow (y \not R x)$ or $R \cap R^{-1} \subseteq \iota_X$.)

Thm. 3.3.12' If R is *antisymmetric*,
 then R is not (always) *reflexive* nor *irreflexive*.

Def. 3.3.12'' R is *asymmetric*, if $\forall x, y \in X, (x, y) \in R \Rightarrow (y, x) \notin R$.
 (or $x R y \Rightarrow y \not R x$) or $R \cap R^{-1} = \emptyset$.

Thm. 3.3.12''' If R is *asymmetric*,
 then R is *irreflexive* and *antisymmetric*.

Def. 3.3.13 R is *transitive*, if $\forall x, y, z \in X: (x, y) \in R \wedge (y, z) \in R \Rightarrow (x, z) \in R$.
 (or $x R y \wedge y R z \Rightarrow x R z$ or $R \cdot R \subseteq R$.)

Def. 3.3.20 Let $R \subseteq X \times X$. Then R is a **partial order**,
if R is **reflexive**, **antisymmetric**, and **transitive**.

Def. 3.3.20' Let $\leq \subseteq X \times X$. Then \leq is a **partial order**. Then
 (X, \leq) is called a **poset** (or **partially ordered sets**).

If $(x \leq y) \vee (y \leq x)$, then x and y are **comparable**.

Otherwise $((x \not\leq y) \wedge (y \not\leq x))$, **incomparable**.

If $\forall x, y \in X$, x and y are comparable \leq is **total** (or **linear**) order.

Exa. 3.3.21 Let $R \subseteq \mathbf{Z}^+ \times \mathbf{Z}^+$. and $(x, y) \in \mathbf{Divides}$, if x divides y . Then
 $(\mathbf{Z}^+, \mathbf{Divides})$ is a **poset**.

(\mathbf{Z}^+, \leq) is a **total order**.

$(2^S, \subseteq)$ is a **poset**.

Hasse Diagram

Let (A, \leq) be a poset. We say a **covers** b , if $a < b \wedge \neg \exists c \in A, a \leq c \wedge c \leq b$.
 (A, covers) is a Hasse Diagram.

Maximal and Minimal Elements

Let (A, \leq) be a poset.

We say a is a **maximal**, if $\neg \exists b \in A, a < b$.

We say a is a **minimal**, if $\neg \exists b \in A, b < a$.

We say a is a **greatest element**, if $\forall b \in A, b \leq a$.

We say a is a **least elements**, if $\forall b \in A, a \leq b$.

Let $S \subseteq A$. $u \in A$ is called **upper bound** of S , if $\forall a \in S, a \leq u$.

$l \in A$ is called **lower bound** of S , if $\forall a \in S, l \leq a$.

$x \in A$ is called the **least upper bound** of S , if $\forall a \in S, a \leq x$,

$x \leq z$, if $\forall z \in \text{upper bound of } S$.

$y \in A$ is called the **greatest lower bound** of S , if $\forall a \in S, y \leq a$,

$z \leq y$, if $\forall z \in \text{lower bound of } S$.

Lattice

A poset (A, \leq) is called **lattice**. If every pair of elements has both a **least upper bound**, and a **greatest lower bound**.

(\mathbf{Z}, \leq) is a lattice.

$\min, \max: \mathbf{Z} \times \mathbf{Z} \rightarrow \mathbf{Z}$

$(\mathbf{Z}^+, |)$ is a lattice.

$\gcd, \text{lcm}: \mathbf{Z}^+ \times \mathbf{Z}^+ \rightarrow \mathbf{Z}^+$

$(2^S, \subseteq)$ is a lattice.

$\cap, \cup: 2^S \times 2^S \rightarrow 2^S$

Def. 3.3.23 Let $R \subseteq X \times Y$, The **inverse** of R , denoted $R^{-1} \subseteq Y \times X$,

$$R^{-1} = \{(y, x) \in Y \times X \mid (x, y) \in R\}$$

Def. 3.3.25 Let $R_1 \subseteq X \times Y$, and $R_2 \subseteq Y \times Z$. The **composition** of R_1 and

R_2 , denoted $R_2 \circ R_1 \subseteq X \times Z$,

$$R_2 \circ R_1 = \{(x, z) \in X \times Z \mid (x, y) \in R_1, (y, z) \in R_2\}$$

Three faces of the relation R from A to B : $R \subseteq A \times B$.

i) R is a set of pairs (subset of Cartesian product)

$$R \subseteq A \times B, (a, b) \in R \text{ where } a \in A \text{ and } b \in B.$$

ii) R is a infix binary boolean (relational) function

$$R: A \times B \rightarrow \{\mathbf{T}, \mathbf{F}\}, a R b \text{ where } a \in A \text{ and } b \in B.$$

R is rectangular boolean matrix. (See Sec. 3.5)

iii) R is a function from A to 2^B .

$$R: A \rightarrow 2^B.$$

$$R(a) = \{b_1, b_2, \dots, b_n\} \text{ where } a \in A \text{ and } \{b_1, b_2, \dots, b_n\} \in 2^B.$$

$$\text{if } 1 \leq \forall i \leq n, (a, b_i) \in R \text{ or } a R b_i \text{ for } n \geq 0,$$

$$\text{if } n=0, R(a) = \{b_1, b_2, \dots, b_n\} = \emptyset.$$

Note that $\forall a \in A, \exists \{b_1, b_2, \dots, b_n\} \in 2^B$. (**unique**)

set valued function.

3.4 Equivalence Relations

Thm 3.4.1 Let S be a partition of a set X . We define a relation on X as

$$R = \{(x, y) \in X \times X \mid x, y \in s, s \in S\}.$$

Then R is **reflexive, symmetric, and transitive**.

proof $\forall x, y, z \in X: \exists s \in S. x, y, z \in s.$

$$\begin{array}{ll} \therefore (x, x) \in R. & \therefore R \text{ is reflexive,} \\ \therefore (x, y) \in R \Rightarrow (y, x) \in R. & \therefore R \text{ is symmetric} \\ \therefore (x, y), (y, x) \in R \Rightarrow (x, z) \in R. & \therefore R \text{ is transitive.} \end{array}$$

Thm 3.4.3 A relation that is **reflexive, symmetric, and transitive** is called an **equivalent relation**.

Exa. $= (\equiv \iota_{\mathbf{N}}) \subseteq \mathbf{N} \times \mathbf{N}$ is an **equivalent relation**

$< \subseteq \mathbf{N} \times \mathbf{N}$ is a **irreflexive total order**

$\leq (\equiv < \cup =) \subseteq \mathbf{N} \times \mathbf{N}$ is a **reflexive total order**

Thm 3.4.8 Let $R \subseteq X \times X$ be an **equivalent** relation on X , let

$[a]_R = \{x \in X \mid a R x\}$ **equivalent class** of X given by the relation R .

Then $S = \{[a]_R \mid a \in X\}$ is a **partion** of X .

proof Since R is **reflexive**, $\forall a \in X: a R a. \therefore a \in [a]_R$.

If $a R b$, then $[a]_R \subseteq [b]_R$, since R is **transitive** and

$[b]_R \subseteq [a]_R$, since R is **symmetric** and **transitive**. \therefore If $a R b$ $[a]_R = [b]_R$.

If $a \not R b$, $[a]_R \cap [b]_R = \emptyset$.

Realtion	$R \subseteq X \times X$	$ R \leq n^2$	$O(n^2)$ where $ X = n$.
Equivalent relation	$S = \{[a]_R \mid a \in X\}$	$ S \leq n$	$O(n)$.

3.5 Matrices of Relations

3.6 Relational Database