

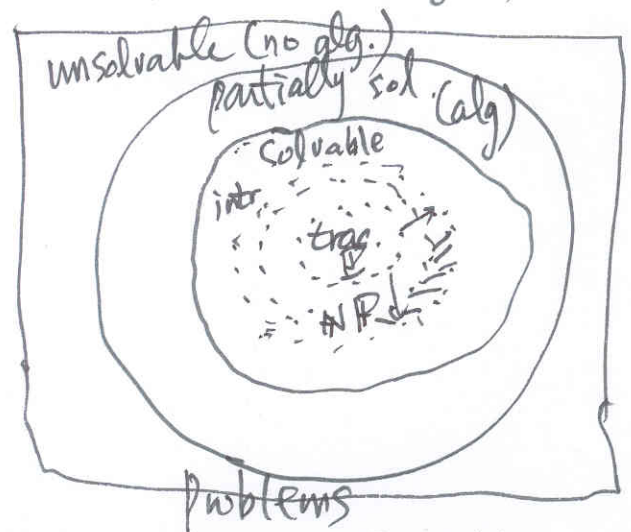
### 3.3 Computational complexity

time complexity  
 space (memory) " (↓)  $\rightarrow$  input data  
 (power consumption)

Worst-case analysis — theoretical  
 Average-case " — practical

linear search  $\frac{n}{2}$   ~~$\frac{n}{2}$~~   $\frac{n}{2}$   $O(n)$   
 ↑ average ↑ worst

binary search  $O(\log_2 n)$

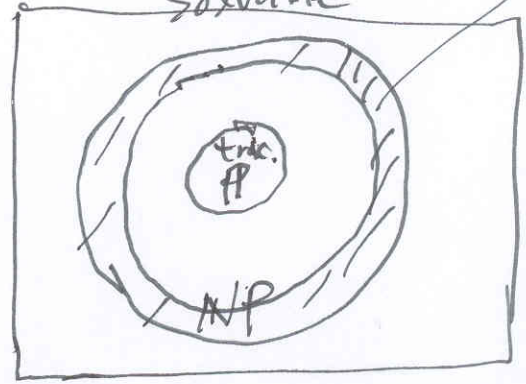


problems

$$P \subseteq NP \begin{cases} P = NP \\ IP \subsetneq NP \end{cases}$$

tractable = class IP  
 $\rightarrow$  poly  $\star$   
 intractable class  
 $\rightarrow$  exp.  $\uparrow$

class NP  
 nondeterministic polynomial.  
 Solvable NP complete



$$3 \quad |\mathbb{N}| = |\mathbb{N}^2| \dots = |\mathbb{N}^k|$$

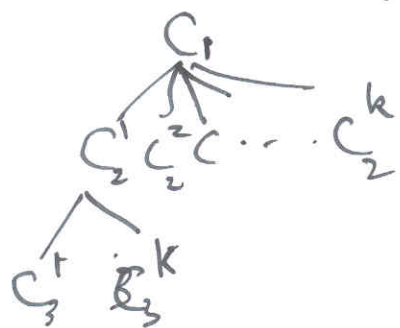
$$|2^{\mathbb{N}}| = \underbrace{(2 \times 2 \times \dots)}_{\text{infinite } \infty}$$

$$\infty^k \underbrace{\infty \times \infty \times \dots \times \infty}_{\text{finite}} \infty$$

### 3.4 The Integer and division on data

$$C_1 \rightarrow \dots \rightarrow C_k$$

$k = O(n)$   
 $\underbrace{\quad}_k$  can be expressed as polynomial ftn of  $n$ .



### Computation

$M, \dots$  memory

$ip \dots$  instruction pointer

$F_{in} \dots$  input data

$F_{out} \dots$  output data

$(M, ip, F_{in}, F_{out})$

↓ Configuration of Computation

$C_0 = (M_0, ip_0, F_{in_0}, F_{out_0})$   
 initial conf.

$$C_0 \rightarrow C_1 \rightarrow C_2 \dots \rightarrow C_n \rightarrow \dots$$

deterministic

$$\forall C_i \exists! C_{i+1}$$

Nondet. otherwise.

$$\exists C_i \begin{matrix} \vdots \\ C_{i+1} \\ \vdots \\ C_{i+k} \end{matrix}$$

### 3.4. Integer and Division

$a, b \in \mathbb{Z}, a \neq 0$ .  $a$  divides  $b$ , if  $\exists c \in \mathbb{Z} \exists b = ac$   
 $a | b$  ( $b \div a \rightarrow$  no remainder)

Ex)  $3 | 6, 3 | 3, 5 \nmid 12, 6 \nmid 3$

$a \equiv b \pmod{a}$  이다.

$a$  is a divisor of  $b$ .

Theorem

i)  $a | b \wedge a | c \rightarrow a | (b+c)$

$a = ds, b = dt \quad s, t \in \mathbb{Z}$

ii)  $a | b \rightarrow a | bc$

$a + b = d(s+t)$

iii)  $a | b \wedge b | c \rightarrow a | c$

Cor. i)  $a | b \wedge a | c \rightarrow a | (mb + mc) \quad \forall m, n \in \mathbb{Z}$

Theorem 2 The division alg.

$\forall a \in \mathbb{Z}, \forall d \in \mathbb{Z}^+, \exists_! q, \exists_! r \in \mathbb{Z}: 0 \leq r < d \exists a = qd + r$

$q = a \text{ div } d, r = a \text{ mod } d$

such that

$\boxed{a \div d = q \dots r}$  초등학급

Modular arithmetic

Def. 3  $a, b \in \mathbb{Z}, m \in \mathbb{Z}^+$

If  $m | (a-b)$  then