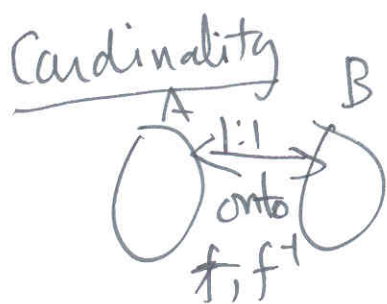
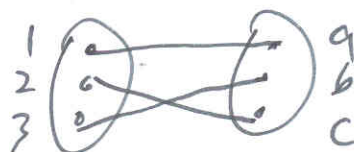


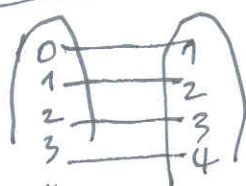
3/16. Uncountable vs algorithms



$$|A| = |B|$$

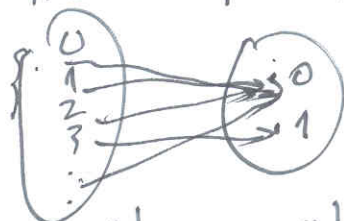


infinite set



Cantor's Diagonal argument

$$f: \mathbb{N} \rightarrow \{0, 1\}$$



(0.10...) infinite binary string
 position 0 1 2 3
 $\{1, \dots\} \subseteq \mathbb{N}$

$$\{0, 1\}^{\mathbb{N}} \cong 2^{\mathbb{N}}$$

infinite binary string

powerset of \mathbb{N}

$$\forall A_0 = (a_{00}, a_{01}, a_{02}, \dots) \quad \forall i, j \in \mathbb{N} \quad a_{ij} \in \{0, 1\}$$

$$A_1 = (a_{10}, a_{11}, a_{12}, \dots)$$

$$A_2 = (a_{20}, a_{21}, a_{22}, \dots)$$

$$A_i = (a_{i0}, a_{i1}, a_{i2}, \dots, a_{ii}, \dots)$$

$$\forall B = (b_0, b_1, b_2, \dots, b_i, \dots)$$

$b_i = 0$ if $a_{ii} = 1$
 $b_i = 1$ if $a_{ii} = 0$
 (otherwise)

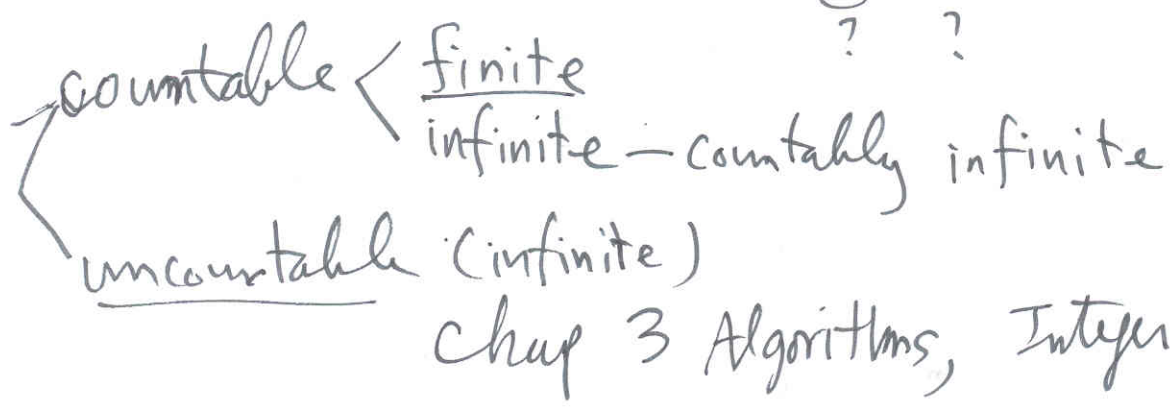
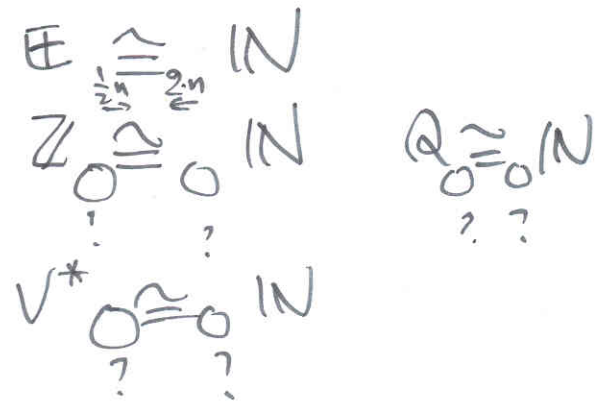
$\mathbb{N} \times \mathbb{N} \dots$ countable $\mathbb{N}^2 \dots$ countable $\mathbb{N}^k \dots$ countable

$$V^* = \bigcup_{i \in \mathbb{N}} V^i = V^0 \cup V^1 \cup V^2 \cup \dots \dots$$

countable

(ϵ), (a, b, ..., z), (aa, ..., zz)

Countable numbering (proof)



3.1 Algorithms.

Dijkstra's mini language

$x := x + 5;$ $x, y := y, x$
 concurrent assignment

$x := x; x := y; y := x$

if B_1 then if B_2 then S_A else S_B fi

dangling else ambiguity

if $B_1 \rightarrow SL_A$
 if $B_2 \rightarrow SL_A$
 fi

if $B_1 \rightarrow SL_1$
 | $B_2 \rightarrow SL_2$
 |
 | $B_n \rightarrow SL_n$
 fi