

2 Basic Structures: Sets, Functions, Sequences, and Sums

2.1 Sets

Definition 1 A set is an unordered collection of objects.

Cantor's naive set theory

Russel's paradox

$$S = \{x \mid x \notin x\}$$

$$x \in S \text{ iff } x \notin x.$$

$$S \in S \text{ iff } S \notin S. \quad \text{contradictory}$$

Definition 2 An object in a set is called **element** or **members** of the set. A set is said to contain its elements.

$a \in A$ “ a is an element of the set A ”

$a \notin A$ “ a is **not** an element of the set A ”

Two way of to define sets

i) To enumerate the elements

$$A = \{a_1, a_2, \dots, a_n\} \quad \text{finite}$$

$$A = \{a_1, a_2, \dots\} \quad \text{infinite}$$

ii) to specify condition with predicate

$$A = \{x \mid P(x)\}$$

$$A = \{x \in U \mid P(x)\} \quad U: \text{universe}(\text{data type})$$

Some important sets in discrete mathematics

$$\mathbf{N} = \{0, 1, 2, 3, \dots\} \quad \text{set of natural numbers.}$$

$$\mathbf{Z} = \{\dots, -2, -1, 0, 1, 2, 3, \dots\} \quad \text{set of integers.}$$

$$\mathbf{Z}^+ = \{1, 2, 3, \dots\} \quad \text{set of positive natural numbers.}$$

$$\mathbf{Q} = \{p/q \mid p \in \mathbf{Z}, q \in \mathbf{Z} \text{ and } q \neq 0\} \quad \text{set of rational numbers.}$$

$$\mathbf{R} \quad \text{set of real numbers.}$$

Definition 3 Two sets are **equal** if and only if they have **same elements**.

$A = B$ “Two set A and B are equal”

$\{\}, \emptyset$ **empty set** “a set that has no elements”

note: $\{\} = \emptyset \neq \{\emptyset\}$.

Definition 4 The set A is said to **subset** of B , if and only if, every elements of A is also an elements of B , and denoted as $A \subseteq B$.

$$A \subseteq B \equiv \forall x(x \in A \rightarrow x \in B)$$

Theorem 1 For any set A ,

$$(i) \emptyset \subseteq S \quad (ii) S \subseteq S.$$

$$(i) \forall S \emptyset \subseteq S \quad (ii) \forall S S \subseteq S$$

For two sets A and B , write $A \subset B$ and say that A is a **proper subset** of B , if and only if, $A \subseteq B$ and (but) $A \neq B$ ($\equiv \neg A = B$).

$$A \subset B \equiv A \subseteq B \wedge A \neq B$$

$$\begin{aligned}
 A = B &\equiv A \subseteq B \wedge B \subseteq A \equiv \forall x(x \in A \rightarrow x \in B) \wedge \forall x(x \in B \rightarrow x \in A) \\
 &\equiv \forall x(x \in A \leftrightarrow x \in B)
 \end{aligned}$$

Definition 5 Let S be a set. If there are exactly n elements in S where n is a nonnegative integer, we say that S is finite set and that n is the **cardinality** of the set S , and denoted as $|S|$.

Definition 6 A set is said to be **finite**, if the cardinality of set is finite. A set is said to be **infinite** if it is not finite.

The Power Set

Definition 7 Given a set S , the power set of S , denoted by $P(S)$, is set of all subsets of S .

$$P(S) \equiv \{A \mid A \subseteq S\}$$

$$|P(S)| = 2^{|S|}.$$

Cartesian Products

Definition 8 The *ordered n -tuple* (a_1, a_2, \dots, a_n) is the *ordered collection*, that has a_1 as its first element, a_2 as its second element, \dots , a_n as its n -th element.

$(a_1, a_2, \dots, a_n) = (b_1, b_2, \dots, b_n)$, iff, $a_i = b_i$ for $i = 1, 2, \dots, n$.

$(a_1, a_2, \dots, a_n) = (b_1, b_2, \dots, b_n) \equiv a_1 = b_1 \wedge a_2 = b_2 \wedge \dots \wedge a_n = b_n$.

We call *ordered pair*, for 2-tuples.

The ordered pairs $(a, b) = (c, d)$, iff, $a = c \wedge b = d$.

Note that $(a, b) \neq (b, a)$

Definition 9 Let A and B be sets. The **Cartesian product**, denoted by $A \times B$, is ...

$$A \times B = \{(a, b) \mid a \in A \wedge b \in B\}.$$

If $R \subseteq A \times B$, the R is called a **relation** from A to B .

$$A \times B \neq B \times A.$$

$$|A \times B| = |A| \times |B|$$

We write $a R b$, if $(a, b) \in R$.

Two aspects of the relation

i) subset of $A \times B$, $(a, b) \in R$.

$$R \subseteq A \times B$$

ii) infix binary boolean operation, $a R b$,

$$R: A \times B \rightarrow \{\mathbf{T}, \mathbf{F}\}$$

Definition 10 The Cartesian product of the sets A_1, A_2, \dots, A_n , denoted by

$$A_1 \times A_2 \times \dots \times A_n = \{(a_1, a_2, \dots, a_n) \mid a_i \in A_i \text{ for } i = 1, 2, \dots, n\}$$

2.2 Set Operations

Definition 1 Union

$$A \cup B = \{x \mid x \in A \vee x \in B\}$$

Definition 2 Intersection

$$A \cap B = \{x \mid x \in A \wedge x \in B\}$$

*Definition 3 Two sets A and B are called **disjoint**, iff, their intersection*

...

$$A \cap B = \emptyset.$$

$$|A \cup B| = |A| + |B| - |A \cap B|$$

principles of set of inclusion-exclusion

Definition 4 Difference

$$A - B = \{x \mid x \in A \wedge x \notin B\}$$

Let universe of discourse be U , U is called **universe**.

Definition 5 Let U be a universe. The complement of the set A , denoted \bar{A} , is called the complement of A with respect to (w.r.t.) U is ...

$$\bar{A} = U - A = \{x \mid x \in U \wedge x \notin A\} = \{x \in U \mid x \notin A\}$$

Four regions in the Venn diagram for two sets A and B .

$$i) A \cap B \quad ii) \bar{A} \cap B \quad iii) A \cap \bar{B} \quad iv) \bar{A} \cap \bar{B}$$

membership table

Four cases for relations on two set in the Venn Diagram

$i) (A \cap \bar{B} = \emptyset) \wedge (\bar{A} \cap B = \emptyset)$	$\equiv A = B,$	equal
$ii) (A \cap \bar{B} = \emptyset) \text{ or } (\bar{A} \cap B = \emptyset)$	$\equiv A \subseteq B \text{ or } B \subseteq A,$	subset
$iii) (A \cap B = \emptyset)$	$\equiv A \cap B = \emptyset,$	disjoint
$iv) \text{ otherwise}$		incomparable

Set Identities

$$A \cup \emptyset = A$$

$$A \cap U = A$$

identity laws

$$A \cup U = U$$

$$A \cap \emptyset = \emptyset$$

domination laws

$$A \cup A = A$$

$$A \cap A = A$$

idempotent laws

$$\overline{\overline{A}} = A$$

double complement law

$$A \cup B = B \cup A$$

$$A \cap B = B \cap A$$

commutative laws

$$A \cup (B \cup C) = (A \cup B) \cup C \quad A \cap (B \cap C) = (A \cap B) \cap C \quad \text{associative ...}$$

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C) \quad A \cap (B \cup C) = (A \cap B) \cup (A \cap C) \quad \text{distributive ...}$$

$$\overline{A \cup B} = \overline{A} \cap \overline{B}$$

$$\overline{A \cap B} = \overline{A} \cup \overline{B}$$

De Morgan's laws

$$A \cup (A \cap B) = A$$

$$A \cap (A \cup B) = A$$

absorption laws

$$A \cup \overline{A} = U$$

$$A \cap \overline{A} = \emptyset$$

complement laws

See Table 6(logical equivalences) of Section 1.2

($\forall, \wedge, \neg, \mathbf{T}, \mathbf{F}$) vs ($\cup, \cap, \overline{}, U, \emptyset$).

boolean algebra
complete lattice

Two prove $A = B$

$$\begin{array}{ll} \text{i) } \forall x(x \in A \rightarrow x \in B) & A \subseteq B, \text{ and} \\ \text{ii) } \forall x(x \in B \rightarrow x \in A) & B \subseteq A. \end{array}$$

Membership table

with n set variables, 2^n membership regions
similar to truth table in logic

Generalized Union and Intersections

Definition 6

$$A_1 \cup A_2 \cup \dots \cup A_n = \bigcup_{i=1}^n A_i$$

Definition 7

$$A_1 \cap A_2 \cap \dots \cap A_n = \bigcap_{i=1}^n A_i$$

Note that \cup and \cap are associative.

2.3 Functions

Definition 1 Let A and B be nonempty sets.

A **function(mapping, transformation)** f from A to B is

an assignment of exactly **one** element of B to **each**(**all**) element of A .

We write $f(a) = b$, if b is the unique element of B assigned to a of A .

We write $f: A \rightarrow B$.

total: to all elements of A (domain)

unique: exactly one elements of B (codomain)

$f: A \rightarrow B$ is a relation from A to B , $f \subseteq A \times B$.

If $f(a)=b$, we write $(a, b) \in f$ or $(a, f(a)) \in f$.

$$f = \{(a, f(a)) \mid \forall a \in A\}$$

$$B^A = \{f: A \rightarrow B\}$$

$$|B^A| = |B|^{|A|}$$

Let $f: A \rightarrow A$, then f is called a function on A .

Definition 2 Let $f: A \rightarrow B$.

A is a **domain** of f , B is a **codomain**(range) of f .

If $f(a)=b$, b is the **image** of a and a is the **preimage** of b .

Two functions f and g are said to be **equal** ... *if $f = g$.(set equivalence)*.

Definition 3 Let $f_1, f_2: A \rightarrow \mathbf{R}$. $f_1+f_2, f_1 f_2: A \rightarrow \mathbf{R}$ is defined by

$$(f_1+f_2)(x) = f_1(x) + f_2(x) \text{ and } f_1 f_2(x) = f_1(x)f_2(x).$$

Definition 4 Let $f: A \rightarrow B$ and $S \subseteq A$. Then the **image** of S under f is

$$\begin{aligned} f(S) &= \{t \in B \mid \exists s \in S (t = f(s))\} \text{ or} \\ &= \{f(s) \in B \mid \forall s \in S\} \text{ for short.} \end{aligned}$$

The range of $f = f(A)$.

One-to-One and Onto function

Definition 5 Let $f: A \rightarrow B$. f is **one-to-one** (1:1) or **injective**, iff

$$\forall a \in A \forall b \in A (a \neq b \rightarrow f(a) \neq f(b)), \quad \text{or logically equivalenty}$$

$$\forall a \in A \forall b \in A (f(a) = f(b) \rightarrow a = b).$$

A injective function is called **injection**.

Definition 6 Let $f: A \rightarrow B$ and (A, \leq) and (B, \leq) are posets (See 8.6),

if $x, y \in A$ and $x < y$, $f(x) \leq f(y)$, f is called **increasing**

$f(x) < f(y)$, f is called **strictly increasing**

$f(x) \geq f(y)$, f is called **decreasing**

$f(x) > f(y)$, f is called **strictly decreasing**

Definition 7 Let $f: A \rightarrow B$. f is **onto** or **surjective**, iff

$$\forall b \in B \exists a \in A (f(a) = b), \quad \text{or} \quad f(A) = B.$$

A surjective function is called **surjection** or **correspondence**.

Definition 8 Let $f: A \rightarrow B$. f is **one-to-one correspondence**, or **bijection**, if f is both **one-to-one**(**injective**) and **onto**(**surjective**).

Theorem 1 Let A and B are sets.

$|A| \leq |B|$, if there is a **injection** $f: A \rightarrow B$.

$|A| \geq |B|$, if there is a **surjection** $f: A \rightarrow B$.

$|A| = |B|$, if there is a **bijection** $f: A \rightarrow B$.

Definition 8.1 Let A and B are sets. We say cardinality of A and B are same, $|A| = |B|$, if there is a **bijection** $f: A \rightarrow B$.

Definition 8.2 Let A and B are sets. and there is a **bijection** $f: A \rightarrow B$. We say that A and B are isomorphic with respect to f , $A \cong_f B$.

Definition 9 Let $f: A \rightarrow B$ and f be one-to-one correspondence. The **inverse function** of f , denoted $f^{-1}: B \rightarrow A$, is defined by

$$f^{-1} = \{(b, a) \mid a \in A, f(a) = b \in B\} \text{ or } f^{-1}(b) = a \text{ when } f(a) = b.$$

Definition 10 Let $g: A \rightarrow B$ and $f: B \rightarrow C$. The **composition** of f and g , denoted by $f \circ g: A \rightarrow C$, is defined by

$$(f \circ g)(a) = f(g(a)) \text{ or } f \circ g = \{(a, c) \mid f(a) = b, g(b) = c\}.$$

Identity function(relation) on A

$$\iota_A = \{(a, a) \mid a \in A\} \quad \text{or } \forall a \in A \iota_A(a) = a.$$

Let $f: A \rightarrow A$. Then

$$f \circ \iota_A = \iota_A \circ f = f.$$

Definition 11 Let $f: A \rightarrow B$. The **graph** of the function f is defined by

$$f = \{(a, b) \mid a \in A, f(a) = b \in B\}$$

Definition 12 The **floor and ceiling function**: $\lfloor \cdot \rfloor \lceil \cdot \rceil: \mathbf{R} \rightarrow \mathbf{Z}$,

$\lfloor x \rfloor = n \in \mathbf{Z}$, n is the **largest integer** such that $n \leq x$.

$\lceil x \rceil = n \in \mathbf{Z}$, n is the **smallest integer** such that $n \geq x$.

2.4 Sequences and Summations

Sequence $\{a_n\}$

$a: \mathbf{N} \rightarrow \mathbf{R}$. We write a_n instead of $a(n)$.

Let $s: \{1, 2, \dots, n\} \rightarrow V$.

We write $s = \text{boy}$ or $s = (b, o, y)$

instead of $s(1) = b, s(2) = o, s(3) = y$.

s is called the **finite string** over V of length n .

V is called the **vocabulary (alphabet)** of string s .

Some Useful Sequences

$n^2, n^3, n^4,$

$2^n, 3^n, n!$

Cardinality

Definition 4 The sets A and B have the same **cardinality**, if and only if, there is a one-to-one correspondence from A to B .

Definition 5 A set is either *finite* or *same cardinality* with \mathbf{Z}^+ (positive integers) is called **countable**. A set that is not countable is called **uncountable**. When an infinite set S is countable, we denote cardinality of S as \aleph_0 (aleph null). $|S| = \aleph_0$.

$$\mathbf{N} \supset \mathbf{Z}^+ \quad \text{but } |\mathbf{N}| = |\mathbf{Z}^+|$$

$$\mathbf{N} \subset \mathbf{Z} \quad \text{but } |\mathbf{N}| = |\mathbf{Z}|$$

$$\mathbf{N} \subset \mathbf{Q} \quad \text{but } |\mathbf{N}| = |\mathbf{Q}|$$

$$\mathbf{N} \subset \mathbf{R} \quad \text{but } |\mathbf{N}| < |\mathbf{R}|$$

Cantor Diagonalization Argument (1879)

Consider $f: \mathbf{N} \rightarrow \{0, 1\}$

f is called **infinite binary string**

Consider cardinality of $\{0, 1\}^{\mathbf{N}} \cong 2^{\mathbf{N}}$.

Assume $|2^{\mathbf{N}}| = \aleph_0$. Then we can *enumerate* binary strings

$$A_0 = (a_{00}, a_{01}, a_{02}, \dots)$$

$$A_1 = (a_{10}, a_{11}, a_{12}, \dots)$$

$$A_2 = (a_{20}, a_{21}, a_{22}, \dots)$$

...

Consider $A = (a_0, a_1, a_2, \dots)$ where $a_n = 0$ if $a_{nn} = 1$; $a_n = 1$, if $a_{nn} = 0$.

$\forall n \in \mathbf{N} A \neq A_n$. But $A \in 2^{\mathbf{N}}$.

\therefore The assumption $|2^{\mathbf{N}}| = \aleph_0$ was wrong.

$\therefore |2^{\mathbf{N}}| > \aleph_0$.

$2^{\mathbf{N}}$ is *uncountable*.

Example 21