

2 Basic Structures: Sets, Functions, Sequences, Sums, and Matrices

2.1 Sets

Def. 1 A set is an **unordered** (distinguished) collection of objects.

An object in a set is called **element** or **member** of the set.

A set is said to **contain** its elements.

$a \in A$ “ a is an **element** of the set A ”

$a \notin A$ “ a is **not** an element of the set A ”

An **ordered** (undistinguished) collection of objects.

ordered pair, triple, quadruple, ..., n -tuple

Two ways to define sets

i) To *enumerate* the elements(원소나열법)

$$A = \{a_1, a_2, \dots, a_n\} \quad \text{finite}$$

$$A = \{a_1, a_2, \dots\} \quad \text{infinite}$$

ii) to specify **condition** with **predicate**(조건제시법)

$$P = \{x \mid p(x)\}$$

$$P = \{x \in U \mid p(x)\} \quad U: \text{universe(domain) of discourse}$$

truth set of predicate $p(x)$

$$P \subseteq U.$$

iii) to write a program(?)

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Def. 2 Two sets are **equal** if and only if they have **same elements**.

$A = B$ “Two set A and B are equal”

$\{\}, \emptyset$ **empty set** “a set that has no elements”

note: $\{\} = \emptyset \neq \{\emptyset\}$.

Def. 3 The set A is said to **subset** of B , if and only if, every elements of A is also an elements of B , and denoted as $A \subseteq B$.

$$A \subseteq B \equiv (\forall x \in A) \Rightarrow (\exists x \in B)$$

For two sets A and B , write $A \subset B$ and say that A is a **proper subset** of B , if and only if, $A \subseteq B$ and (but) $A \neq B$ ($\Leftrightarrow \neg(A = B)$).

$$A \subset B \equiv (A \subseteq B) \wedge (A \neq B)$$

New Definition for Equality of sets

Def. 2.1 Two sets A and B are **equal** if and only if A is a subset of B and B is a subset of A (vice versa).

$$\begin{aligned} A = B, & \quad \equiv (A \subseteq B) \wedge (B \subseteq A) \\ & \quad \equiv (\forall x \in A) \Rightarrow (x \in B) \wedge (\forall x \in B) \Rightarrow (x \in A) \\ & \quad \equiv (\forall x \in A) \Leftrightarrow (x \in B) \end{aligned}$$

Two prove $A = B$

$$\begin{array}{ll} i) & \forall x \in A \Rightarrow x \in B) & A \subseteq B, \text{ and} \\ ii) & \forall x \in B \Rightarrow x \in A) & B \subseteq A. \end{array}$$

Theorem 1 For any set S ,

- (i) $\emptyset \subseteq S$ (ii) $S \subseteq S$.
 (i) $\forall S(\emptyset \subseteq S)$ (ii) $\forall S(S \subseteq S)$

Def. 4 Let S be a set. If there are **exactly** n elements in S where n is a non-negative integer, we say that S is **finite** set and that n is the **cardinality** of the set S , and denoted as $|S|$ (or $\#(S)$).

Def. 5 A set is said to be **finite**, if the cardinality of set is finite. A set is said to be **infinite** if it is not finite.

The Power Set

Def. 6 Given a set S , the power set of S , denoted by $P(S)$ or 2^S , is set of all subsets of S .

$$P(S) = 2^S \equiv \{A \mid A \subseteq S\}$$

$$|P(S)| = |2^S| = 2^{|S|}.$$

Cartesian Products

Def. 8 Let A and B be sets. The **Cartesian product** of set A and B , denoted by $A \times B$, is set of **all ordered pairs** (a, b) , where ...

$(a, b) \in A \times B$, $A \times B = \{(a, b) \mid a \in A, b \in B\}$., means \wedge in 조건 제시법

$$|A \times B| = |A| \times |B|$$

$$A \times B \neq B \times A$$

The ordered pairs $(a, b) = (c, d)$, iff, $(a = c) \wedge (b = d)$.

Note that $(a, b) \neq (b, a)$ but $\{a, b\} = \{b, a\}$.

Def. 9 The Cartesian product of the sets A_1, A_2, \dots, A_n , denoted by

$$A_1 \times A_2 \times \dots \times A_n = \{(a_1, a_2, \dots, a_n) \mid a_i \in A_i \text{ for } i = 1, 2, \dots, n\}$$

Def. 7 The **ordered n -tuple** (a_1, a_2, \dots, a_n) is ...

$$(a_1, a_2, \dots, a_n) \in A_1 \times A_2 \times \dots \times A_n.$$

2.2 Set Operations

Def. 1 Union

$$A \cup B = \{x \mid x \in A \vee x \in B\} = \{x \mid x \in A \text{ or } x \in B\}.$$

Def. 2 Intersection

$$A \cap B = \{x \mid x \in A \wedge x \in B\} = \{x \mid x \in A, x \in B\}.$$

Def. 3 Two sets A and B are called **disjoint**, iff, their intersection ...

$$A \cap B = \emptyset.$$

$$|A \cup B| = |A| + |B| - |A \cap B|$$

principles of set of inclusion-exclusion

Def. 4 Difference

$$A - B = \{x \mid x \in A, x \notin B\} = \{x \mid x \in A \wedge x \notin B\}.$$

Let domain of discourse for sets be U , U is called **universe** of the set.

Def. 5 Let U be a **universe**. The **complement** of the set A , denoted \bar{A} , is called the complement of A with respect to (w.r.t.) U is ...

$$\bar{A} = U - A = \{x \mid x \in U, x \notin A\} = \{x \in U \mid x \notin A\}$$

Set Identities**Table 1 in p. 132****1. Identity laws**

$$A \cup \emptyset = A$$

$$A \cap U = A$$

2. Domination laws

$$A \cup U = U$$

$$A \cap \emptyset = \emptyset$$

3. Idempotent laws

$$A \cup A = A$$

$$A \cap A = A$$

4. Double Complement law

$$\overline{\overline{A}} = A$$

5. Commutative laws

$$A \cup B = B \cup A$$

$$A \cap B = B \cap A$$

6. Associate laws

$$A \cup (B \cup C) = (A \cup B) \cup C$$

$$A \cap (B \cap C) = (A \cap B) \cap C$$

7. Distributive laws

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

8. De Morgan's laws

$$\overline{A \cup B} = \overline{A} \cap \overline{B}$$

$$\overline{A \cap B} = \overline{A} \cup \overline{B}$$

9. Absorption laws

$$A \cup (A \cap B) = A$$

$$A \cap (A \cup B) = A$$

10. Complement laws

$$A \cup \overline{A} = U$$

$$A \cap \overline{A} = \emptyset$$

11. U/\emptyset laws

$$\overline{U} = \emptyset$$

$$\overline{\emptyset} = U$$

Compare with the Table 6(logical equivalence) of Section 1.2 p. 25

$(\{\mathbf{T}, \mathbf{F}\}, \neg, \vee, \wedge)$ vs $(\{U, \emptyset\}, \bar{}, \cup, \cap)$

propositional logic and set algebra are **isomorphic**

$(\{\mathbf{T}, \mathbf{F}\}, \Rightarrow)$ and complete lattice

Four regions in the Venn diagram for two sets A and B .

i) $A \cap B$ ii) $\bar{A} \cap B$ iii) $A \cap \bar{B}$ iv) $\bar{A} \cap \bar{B}$

Compare with the truth table in the **logic**

2^n regions in the Ven diagram for n sets

Four cases for relations on two set in the Venn Diagram

i) $(A \cap \bar{B} = \emptyset) \wedge (\bar{A} \cap B = \emptyset)$	$\equiv A = B,$	equal
ii) $(A \cap \bar{B} = \emptyset)$ or $(\bar{A} \cap B = \emptyset)$	$\equiv A \subseteq B$ or $B \subseteq A,$	subset
iii) $(A \cap B = \emptyset)$	$\equiv A \cap B = \emptyset,$	disjoint
iv) otherwise		incomparable

Generalized Union and Intersections

Definition 6 Generalized Union

$$A_1 \cup A_2 \cup \dots \cup A_n = \bigcup_{i=1}^n A_i = \bigcup_{i \in N_n} A_i \text{ where } N_n = \{1, 2, \dots, n\}.$$

Definition 7 Generalized Intersection

$$A_1 \cap A_2 \cap \dots \cap A_n = \bigcap_{i=1}^n A_i = \bigcap_{i \in N_n} A_i \text{ where } N_n = \{1, 2, \dots, n\}.$$

Note that \cup and \cap are **associative** and n -ary operator.

2.3 Functions

Def. 1 Let A and B be nonempty sets.

A **function**(mapping, transformation) f from A to B is

an assignment of **exactly one** element of B to **each**(all) element of A .

We write $f(a) = b$, if b is the **unique** element of B assigned to a of A .

We write $f: A \rightarrow B$, if f is a function from A to B .

total: for **all** elements of A (**domain**)

unique: exactly **one** elements of B (**codomain**)

$\forall a \in A: \exists^1 b \in B . \exists. f(a) = b.$

$|A| <^? |B|.$

A **function** can be considered as a **set of pairs**!(Remember the **graph**)

$$f = \{(a, f(a)) \mid a \in A, f(a) \in B\}$$

Let $f: A \rightarrow A$, then f is called a function **on** A .

Def. 2 Let $f: A \rightarrow B$.

A is a **domain** of f , B is a **codomain(range)** of f .

If $f(a)=b$, b is the **image** of a and a is the **preimage** of b .

Two functions f and g are said to be **equal**, if $f = g$.(set equivalence).

Def. 3 Let $f_1, f_2: A \rightarrow \mathbf{R}$. $f_1+f_2, f_1 f_2: A \rightarrow \mathbf{R}$ is defined by

$$(f_1+f_2)(x) = f_1(x) + f_2(x) \text{ and } f_1 f_2(x) = f_1(x)f_2(x).$$

Def. 4 Let $f: A \rightarrow B$ and $S \subseteq A$. Then the imange of S under f is

$$\begin{aligned} f(S) &= \{t \in B \mid \exists s \in S, t = f(a)\} \text{ or} \\ &= \{f(s) \in B \mid \forall s \in S\} \text{ for short.} \\ &\subseteq B. \end{aligned}$$

The range of f is $f(A) \subseteq B$.

How many different functions $f: A \rightarrow B$?

One-to-One and Onto function

Def. 5 Let $f: A \rightarrow B$. f is **one-to-one**(1:1) or **injective**, iff

$$\forall a \in A \quad \forall b \in A, [(f(a)=f(b)) \rightarrow (a=b)] \quad \text{or logically equivalent}$$

$$\forall a \in A \quad \forall b \in A, [(a \neq b) \rightarrow (f(a) \neq f(b))].$$

A injective function is called **injection**.

$$|A| \leq |B|.$$

Def. 6 Let $f: A \rightarrow B$ and (A, \leq) and (B, \leq) are **posets**(See 9.6),

if $x, y \in A$ and $x < y$, $f(x) \leq f(y)$, f is called **increasing** and

$f(x) < f(y)$, f is called **strictly increasing**

$f(x) \geq f(y)$, f is called **decreasing**

$f(x) > f(y)$, f is called **strictly decreasing**

Def. 7 Let $f: A \rightarrow B$. f is **onto** or **surjective**, iff

$$\forall b \in B, \exists a \in A . \exists . f(a) = b, \quad \text{or} \quad f(A) = B.$$

A surjective function is called **surjection** or **correspondence**.

$$|A| \geq |B|.$$

Def. 8 Let $f: A \rightarrow B$. f is **one-to-one correspondence** or **one-to-one onto** or **bijective**, if f is both **one-to-one (injective)** and **onto (surjective)**.

$$|A| = |B|.$$

Def. 9 Let $f: A \rightarrow B$ and f be **one-to-one** and **onto (bijective)**.

The **inverse** of f also is a **function**, denoted $f^{-1}: B \rightarrow A$, is defined by

$$f^{-1} = \{(b, a) \mid a \in A, f(a) = b \in B\} \text{ or}$$

$$f^{-1}(b) = a \text{ when } f(a) = b.$$

Def. 10 Let $g: A \rightarrow B$ and $f: B \rightarrow C$. The composition of f and g , denoted by $f \circ g: A \rightarrow C$, is defined by

$$(f \circ g)(a) = f(g(a)) \text{ or } f \circ g = \{(a, c) \mid f(a) = b, g(b) = c\}.$$

Def. 11 Let $f: A \rightarrow B$. The **graph** of the function f is defined by

$$f = \{(a, b) \in A \times B \mid a \in A, f(a) = b \in B\}$$

Def. 12 The **floor** and **ceiling** function: $\lfloor \cdot \rfloor \lceil \cdot \rceil : \mathbf{R} \rightarrow \mathbf{Z}$,

$$\lfloor x \rfloor = n \in \mathbf{Z}, n \text{ is the largest integer such that } n \leq x. (\text{floor})$$

$$\lceil x \rceil = n \in \mathbf{Z}, n \text{ is the smallest integer such that } n \geq x. (\text{ceiling})$$

Def. 13 partial function $f: A \mapsto B$.

$\forall a \in A: \exists b \in B . \exists f(a) = b$. in **Definition 1**. But

$\exists a \in A: . \exists f(a)$ is undefined.

~~total~~ and uniqueness

Identity function(relation) on A

Let $f: A \rightarrow A$. Then

$$\iota_A = \{(a, a) \mid a \in A\}$$

$$\text{or } \forall a \in A \iota_A(a) = a.$$

$$f \circ \iota_A = \iota_A \circ f = f.$$

Identity function w.r.t function composition.

2.4 Sequences and Summations

Def. 1 sequence (수열) $\{a_n\}$

$a: \mathbf{N} \rightarrow \mathbf{R}$. We write a_n instead of $a(n)$.

Def. 2 Geometric sequence(*progression*; 등비수열)

$$\{g_n\} = g(n) = ar^n.$$

Def. 3 Arithmetic sequence(*progression*; 등차수열)

$$\{b_n\} = b(n) = a + nd.$$

Ex. 4 string: Let $s: \{1, 2, \dots, n\} \rightarrow V = \{a, b, \dots, z\}$.

We write $s = \text{boy}$ or $s = (b, o, y)$

instead of $s(1) = b, s(2) = o, s(3) = y$.

V is called the **vocabulary**(*alphabet*) of string s .

s is the **string**(*vocabulary sequence* 문자열) over V of length n .

Recurrence Relation (점화식)

Def. 4 Let $a: N \rightarrow \mathbf{R}$ and $f: \mathbf{R}^n \rightarrow \mathbf{R}$. We define $a: N \rightarrow \mathbf{R}$ by
 $a_n = f(a_0, a_1, \dots, a_{n-1})$ with definitions of a_0, a_1, \dots , and a_{n-1} .
 $a(n) = f(a(0), a(1), \dots, a(n-1))$ with definitions of $a(0), \dots, a(n-1)$.

Def. 5 Fibonacci sequence

$$f_n = f_{n-1} + f_{n-2} \text{ with } f_0 = 0 \text{ and } f_1 = 1$$

Special Integer Sequences

polynomial sequences n, n^2, n^3, n^4, \dots

exponential sequences $2^n, 3^n, \dots, n!$

See Table 1 in p. 162

Summations

See Theorem 1 and Table 2 in p. 166

2.5 Cardinality of Sets

Def. 1 A set S is **finite** with cardinality $n \in \mathbf{N}$, written $|A| = n$, if there is a bijection from the set $\{0, 1, \dots, n-1\}$ to S .
A set is **infinite** otherwise.

Def. 2 The sets A and B have the same **cardinality**, if and only if, there is a **bijection** (one-to-one correspondence) from A to B .

Thm. 3.0 Let A and B be sets.

$|A| \leq |B|$, if there is a **injection** (one-to-one) $f: A \rightarrow B$.

$|A| \geq |B|$, if there is a **surjection** (onto): $A \rightarrow B$.

$|A| = |B|$, if there is a **bijection** $f: (one-to-one \wedge onto) A \leftrightarrow B$.

Def. 3.1 Let A and B be sets. We say the **cardinalities** of A and B are same, $|A| = |B|$, if there is a **bijection** $f: A \leftrightarrow B$.

Extended Set Equivalence

We say that two sets A and B are **isomorphic** with respect to f , written $A \cong_f B$ or $A \cong B$ for short.

If f is a **bijection** from A to B and vice versa. $f: A \leftrightarrow B$.

$$\forall a \in A \exists! f(a) \in B \text{ and } \forall b \in B \exists! f^{-1}(b) \in A.$$

We can identify B from A and f , and identify A from B and f^{-1} (vice versa)

Set Isomorphism includes Set Equivalence

$\{0, 1, \dots, n-1\} \cong_f \{1, 2, \dots, n\}$ What is f ?

$\{0, 1, \dots, n-1\} = \{0, 1, \dots, n-1\}$ and $\{0, 1, \dots, n-1\} \cong_f \{0, 1, \dots, n-1\}$ f ?

Countable Sets

Def. 4 A set is either **finite** or **same cardinality** with \mathbf{Z}^+ (positive integers) is called **countable**. A set that is not countable is called **uncountable**. When an infinite set S is countable, we denote cardinality of S as \aleph_0 (aleph null). $|S| = \aleph_0$.

$$\mathbf{N} \supset \mathbf{N}^+ \quad \text{but } |\mathbf{N}| = |\mathbf{N}^+| = \aleph_0.$$

$$\mathbf{N} \cong_f \mathbf{N}^+ \quad \text{What is the natural number, } \mathbf{N} \text{ or } \mathbf{N}^+ ?$$

$$\mathbf{I} \supset \mathbf{N} \quad \text{but } |\mathbf{I}| = |\mathbf{N}| = \aleph_0. \quad \text{What is the **bijection**?$$

$$\mathbf{Q} \supset \mathbf{N} \quad \text{but } |\mathbf{Q}| = |\mathbf{N}| = \aleph_0. \quad \text{See example 4}$$

But!

$$\mathbf{R} \supset \mathbf{N} \quad \text{but } |\mathbf{R}| > |\mathbf{N}| = \aleph_0. \quad \text{See example 5}$$

Uncountable Sets

Cantor Diagonal Argument(1874, 1891)

Consider $f: \mathbf{N} \rightarrow \{0, 1\}$ or $f = \{0, 1\}^{\mathbf{N}}$.

f is an *infinite binary string*

Consider the cardinality of $\{0, 1\}^{\mathbf{N}} \cong 2^{\mathbf{N}}$.

Power set of natural numbers

Assume $|2^{\mathbf{N}}| = |\mathbf{N}|$. Then we can **enumerate**(번호매김 ; countable)
binary strings B_i for $i \in \mathbf{N}$.

$$B_0 = (b_{00}, b_{01}, \dots, b_{0n}, \dots)$$

$$B_1 = (b_{10}, b_{11}, \dots, b_{1m}, \dots)$$

...

$$B_n = (b_{n0}, b_{n1}, \dots, b_{nn}, \dots)$$

...

Define **complement** of a binary string $B = (b_0, b_1, \dots, b_n, \dots)$,

$$\bar{B} = (\bar{b}_0, \bar{b}_1, \dots, \bar{b}_n, \dots) \text{ where for } i \in \mathbf{N},$$

$$\bar{b}_i = 0, \text{ if } b_i = 1 \text{ and } \bar{b}_i = 1, \text{ if } b_i = 0.$$

Consider an infinite **diagonal** binary string $B_d = (b_{00}, b_{11}, \dots, b_{nn}, \dots)$
and its **complement** $\bar{B}_d = (\bar{b}_{00}, \bar{b}_{11}, \dots, \bar{b}_{nn}, \dots)$

$\forall n \in \mathbf{N}, \bar{B}_d \neq B_n$. But $\bar{B}_d \in 2^{\mathbf{N}}$. (\bar{B}_d is an infinite binary string)

$$\therefore |2^{\mathbf{N}}| > |\mathbf{N}|.$$

\therefore The assumption $|2^{\mathbf{N}}| = |\mathbf{N}|$ was **wrong**.

$2^{\mathbf{N}}$ is **uncountable**.

Note that core of the proof is

complement of diagonal (**denial** of self recursion)

in infinite string

Halting problem

*Is there a program that reads program as a data, and
decides whether the program with the data will **halt or not**?*

In Chap. 3.1

Some similar examples in the world

*A barber who shave everybody who can **not** shave himself.*

*Shall the barber shave **himself**?*

*An adjective is heterological, if the adjective does **not** possess
the property it describes. (monosyllabic, polysyllabic)*

*Is the adjective “**heterological**” **heterological**?*

How about “homological”?

*There is a sign that “It is written by **me**(**liar**)”.*

*Did you(**liar**) write it?*

Contradiction

denial of self recursion!

Russel's paradox(1901, 1911)

Assume a set R is defined as $R = \{x \mid x \notin R\}$.

Consider x as R itself, Then $R \in R$ iff $R \notin R$. **contradictory**

$\therefore R$ is undefined!

Self denial is a contradiction

How about the set?

$$\bar{R} = \{x \mid x \in R\}$$

Gödel's Incompleteness Theorem(1931)

Hilbert's 2nd problem(23 problems; 1929)

Is there a

Any logic cannot be both consistent and complete.

proof

$G(p: \text{Theorem}) \equiv$ Gödel number of the theorem p

$F(x: \text{Gödel number}) \equiv$

if $\exists p: \text{Theorem}(\text{Proof})$ for $x \rightarrow \mathbf{T}$

| $\nexists p: \text{Theorem}(\text{Proof})$ for $x \rightarrow \mathbf{F}$ **fi**

$\overline{Bew}(y: \text{Gödel number}) \equiv$

if $\exists p: F(G(p)) \rightarrow \mathbf{T}$ **|** $\nexists p: F(G(p)) \rightarrow \mathbf{F}$ **fi**

$\overline{\overline{Bew}}(y: \text{Gödel number}) \equiv$

if $\exists p: F(G(p)) \rightarrow \mathbf{F}$ **|** $\nexists p: F(G(p)) \rightarrow \mathbf{T}$ **fi**

$\overline{Bew}(G(\overline{Bew})) \Leftrightarrow \exists p: F(G(\overline{Bew}))$

$\Leftrightarrow \nexists p: F(G(\overline{Bew}))$

Contradiction!

Thm. 2 If A and B are **countable** sets, then $A \cup B$ is also **countable**.

proof Without loss of generality, we can assume A and B are **disjoint**.

If there are not, we can replace B by $B - A$ because

$$A \cap (B - A) = \emptyset \text{ and } A \cup (B - A) = A \cup B. \text{ (Venn diagram)}$$

case i) A and B are both **finite**. $|A| + |B|$ is finite and **countable**.

case ii) A **countable** and B is **finite**.

$$A \cup B = b_1, b_2, \dots, b_m, a_1, a_2, \dots, a_n, \dots \quad \text{is countable.}$$

case iii) Both A and B are **countable**.

$$A \cup B = a_1, b_1, a_2, b_2, \dots, a_n, b_m, \dots \quad \text{is countable.}$$

Theorem 3 Let A and B are sets with $|A| \leq |B|$ and $|B| \leq |A|$. Then $|A| = |B|$.

In other words,

$$\exists f: A \rightarrow B \text{ is a 1-1 and } (\exists g: B \rightarrow A \text{ is}) \text{ 1-1.}$$

Definition 5 Let S be a set. Then $f: \mathbb{N} \rightarrow S$

2.6 Matrices