

4/6(4) Recursion .

Recursion (Induction)

Basis $P(n_0)$.

recus: $P(k) \Rightarrow P(k+1)$
 $\forall k \geq n_0$.

$P(n_0)$
 $P(n_0+1)$ — C-basis
 \vdots
 $P(n_0+C-1)$

$k \in \mathbb{N}, k \leq C$

Recursion은

~~recursion~~
 귀납의 $\frac{1}{2}$ 은 $\frac{1}{2}$ 은 $\frac{1}{2}$ 은 $\frac{1}{2}$ 은
 (induction) (Deduction)

Ex) $a_0 = b$

$a_n = a \cdot a_{n-1}$ ($n \geq 1$)
 $n > 0$... 점화식 (recurrence relation) \rightarrow chap. 8.

$f_0, f_1 := 0, 1$;

$f_n = f_{n-1} + f_{n-2}$ ($n \geq 2$) ... 2차 점화식

k-차 점화식

의 바깥 (closed form) 구하기 $a_n = f(a_{n-1}, \dots, a_{n-k}, f(0), f(1), \dots, f(k))$

~~$(a_n = a^n)$~~ - 귀납 ... plug-and-chug
 대입하면 - 고쳐쓰고

$a_0 = b$
 $a_1 = a \cdot a_0 = a \cdot b = a \cdot b$
 $a^2 = a \cdot a_1 = a \cdot ab = a^2 \cdot b$
 $a^n = ba^n$

(증명) $a^n = ba^n$ 이다 $n \geq 0$
 basis $a_0 = b \cdot a^0 = b$ $a = b$ (by def)
 rec. $a_k = ba^k$ 이다 가정

$a_{k+1} = a \cdot a_k = a \cdot ba^k = b \cdot a^{k+1}$
 점화식에 I.H. 모양이 맞기.

~~$f_0, f_1 = 0, 1;$~~
 function fibo($n \in \mathbb{N}$) $\in \mathbb{N}$. $\left[\begin{array}{l} f_0 = 0, f_1 = 1 \\ f_n = f_{n-1} + f_{n-2} \quad n \geq 2 \end{array} \right.$

if $n=0 \rightarrow$ return 0

$\square n=1 \rightarrow$ return 1

$\square n \geq 2 \rightarrow$ ~~fibo($n-1$) + fibo($n-2$); return fibo~~

fi
~~return~~

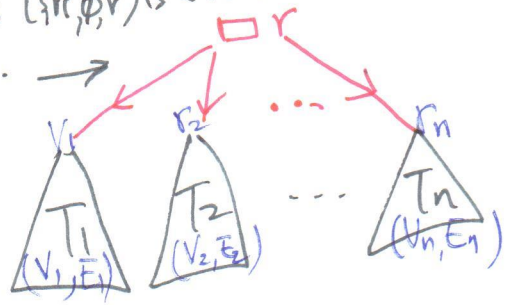
E.W. Dijkstra

Code for machine vs human-being!
 (mathematics (?))

P.L.

rooted Tree
 basis $(\text{root}, \emptyset, r)$ is a r.t.

ind. \rightarrow



한글이 잘쓰기

1. 번역의 탄생
 2. 부한대, 문장교장기
- 한글이

<http://speller.cs.pusan.ac.kr/>